

CPSC 4240/5240: Parallel Programming Techniques

Lecture 10: Concurrent Data Structures Continued

Lecturer: Quanquan C. Liu
Spring 2026

Homework 2 is Out

- Due **next Thursday, late day Monday, March 2nd**
- See Files/Homework 2

Concurrent Linked List: Challenges

- Consider the following linked list and scenario
- Thread 1 wants to **remove C**
- Thread 2 wants to **add D before E**

**What could
possibly
happen here?**



Concurrent Linked List: Challenges

- Consider the following scenario
- Thread 1 wants to **remove C**
- Thread 2 wants to **add D before**

Example execution:

- 1) Thread 1 removes pointer to E
and pauses

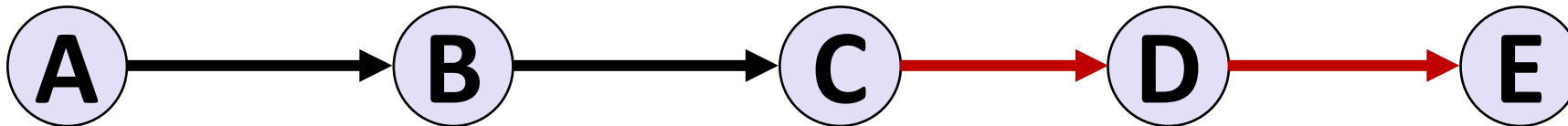


Concurrent Linked List: Challenges

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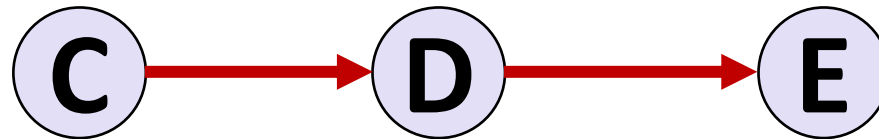


Concurrent Linked Lists

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Concurrent Linked Lists

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and pauses
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Lots of issues!

- Dangling pointer from D from empty memory address
 - Two pointers into E
 - No way to get to D

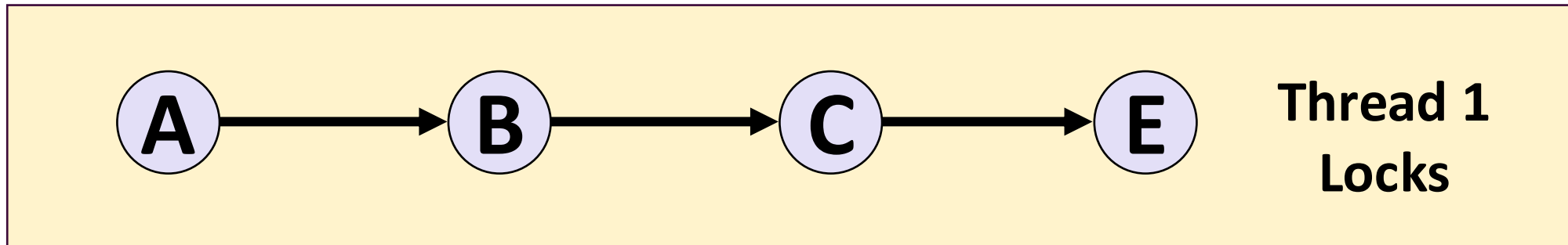
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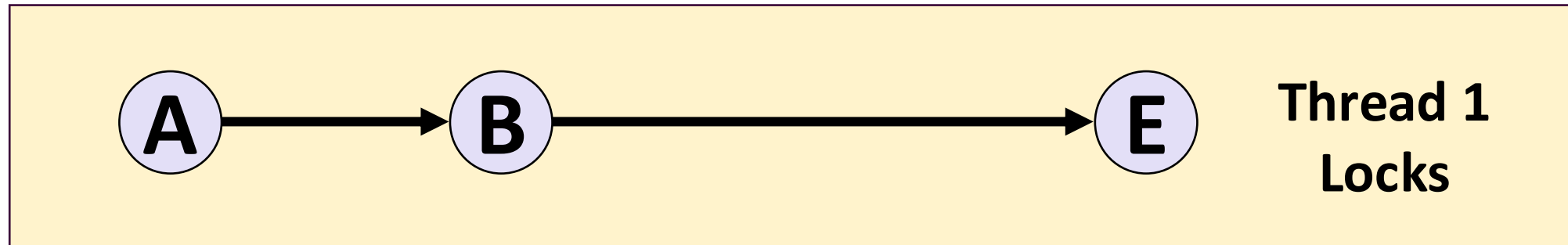
Concurrent Linked List using Locks

- Can just lock the entire linked list when you perform any operation
 - **Course-grained locking**



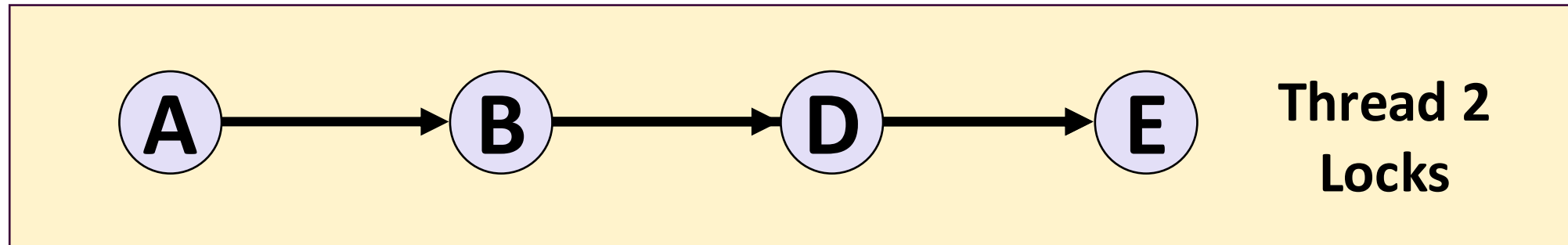
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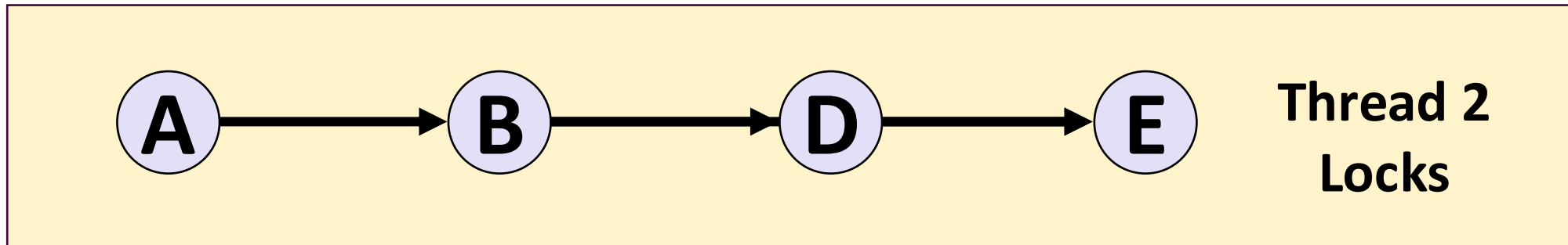
Concurrent Linked List using Locks

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Concurrent Linked List using Locks

- Can just lock the entire linked list when you perform any operation
 - **Course-grained locking**
 - **Not good because too sequential, not using multicore processors at all**



Concurrent Linked List using Locks

- Fine-grained locking: what can you do?

Concurrent Linked List using Locks

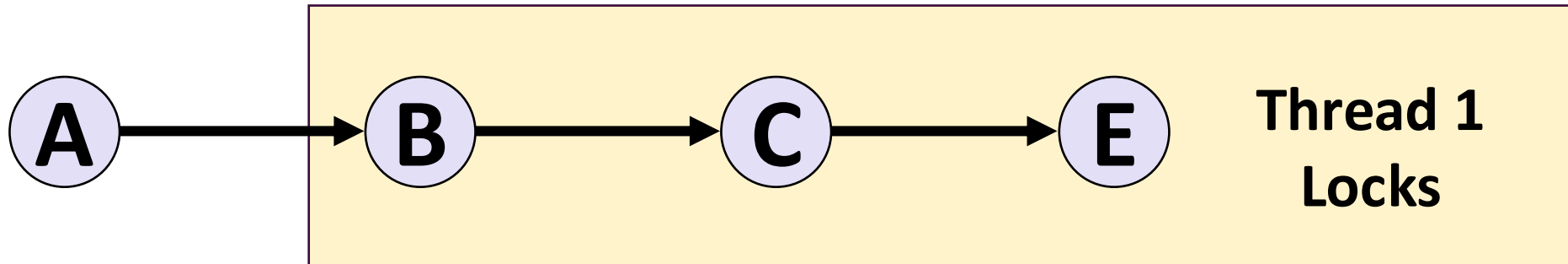
- **Fine-grained locking**

- Lock only the nodes that are affected by each operation
- E.g. only lock the neighbors!

Concurrent Linked List using Locks

- **Fine-grained locking**

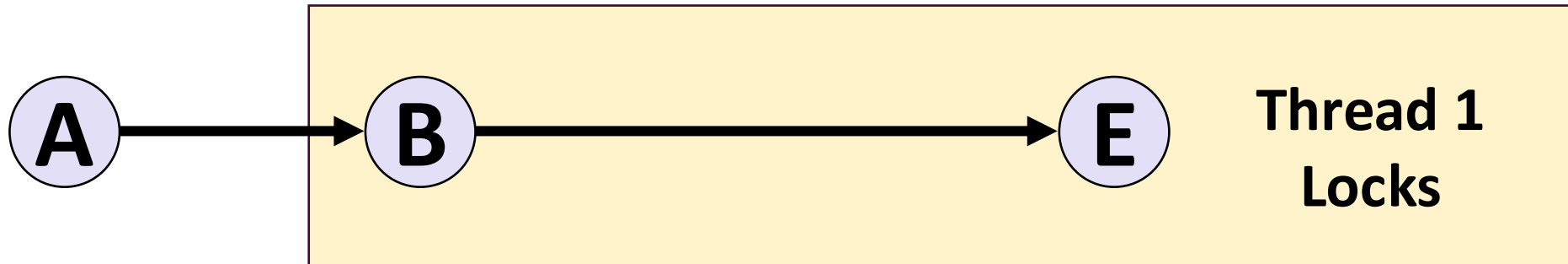
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Concurrent Linked List using Locks

- **Fine-grained locking**

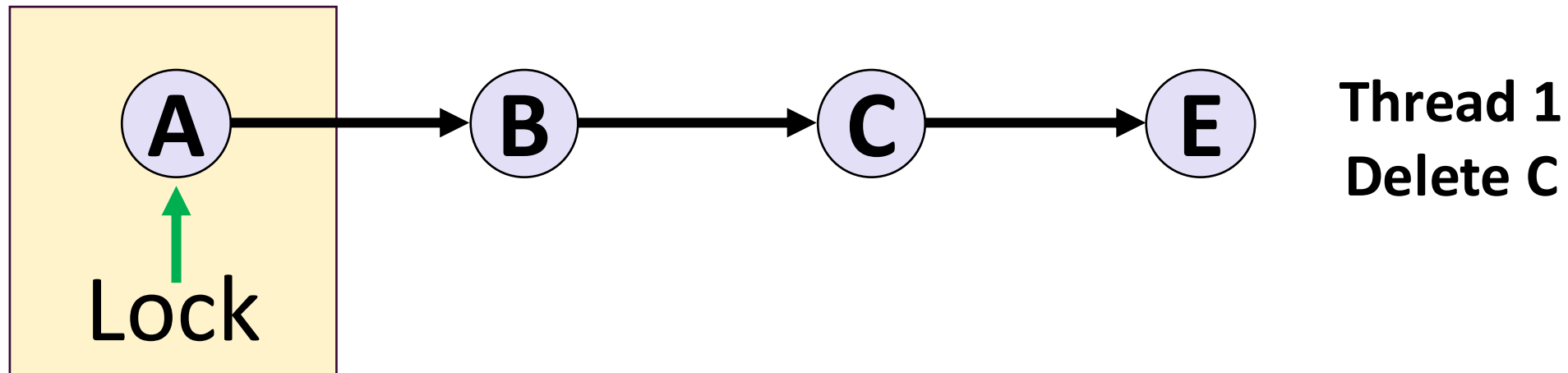
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Concurrent Linked List using Locks

- **Deletion example step by step:**

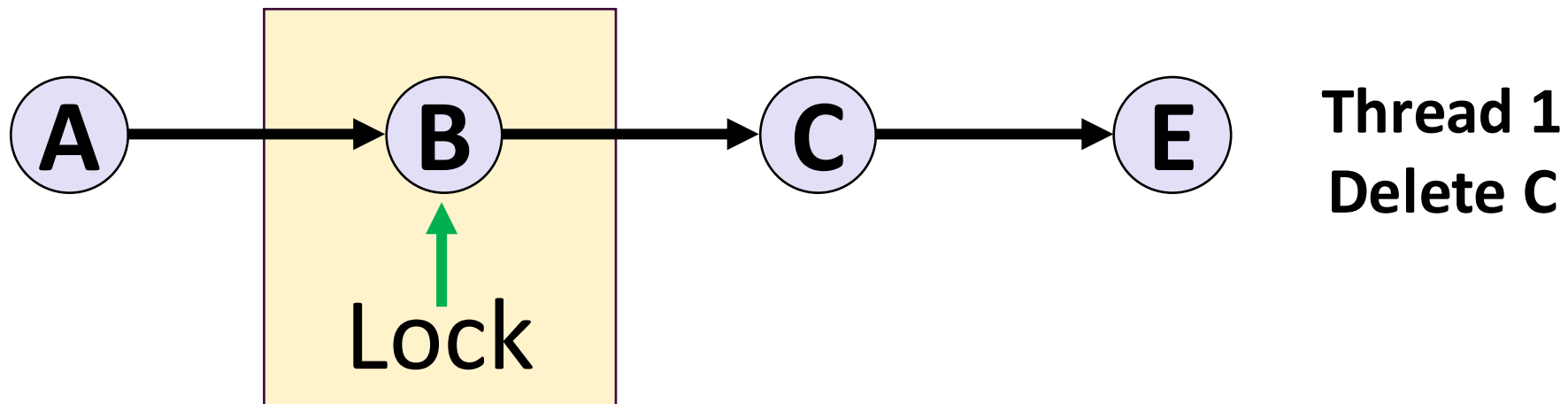
- First, start from the head node and find the element in the list
- What do we need to do as we are trying to find the element in the list? What can go wrong?
- We need to lock each node as we traverse because the node could have been changed while we are traversing



Concurrent Linked List using Locks

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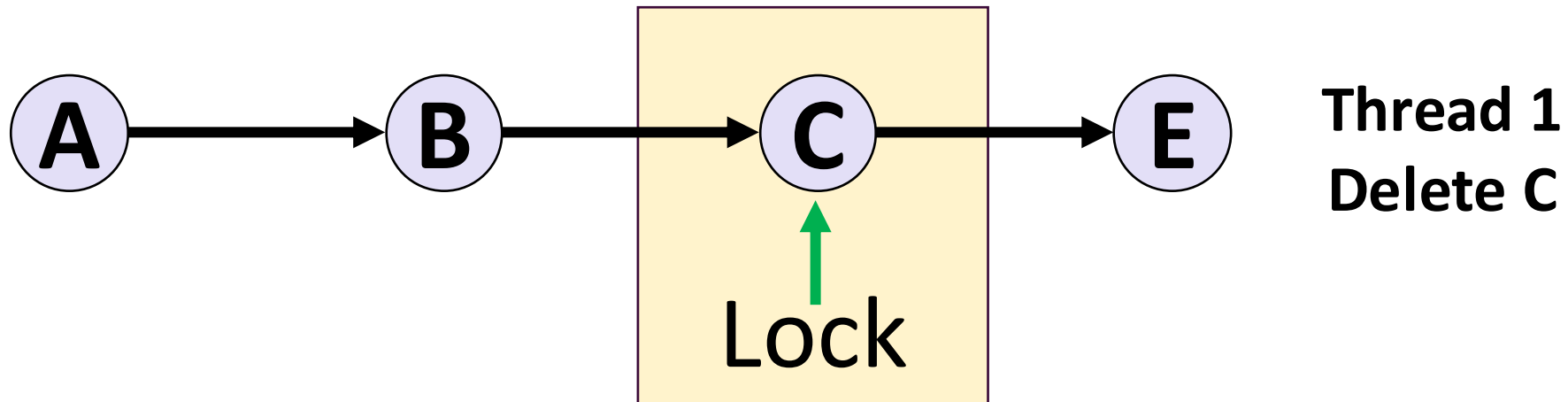
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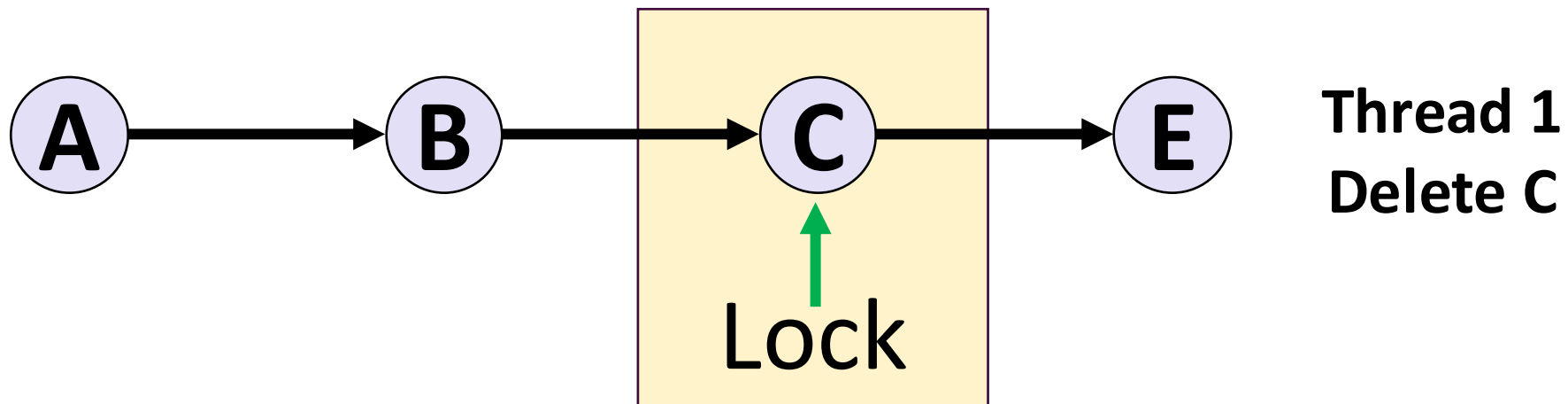
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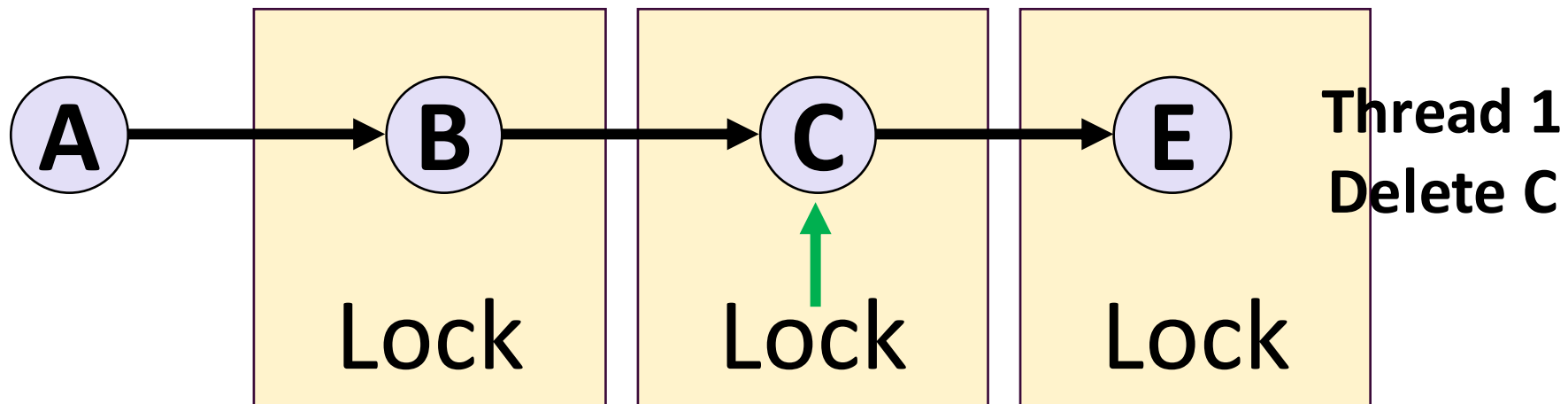
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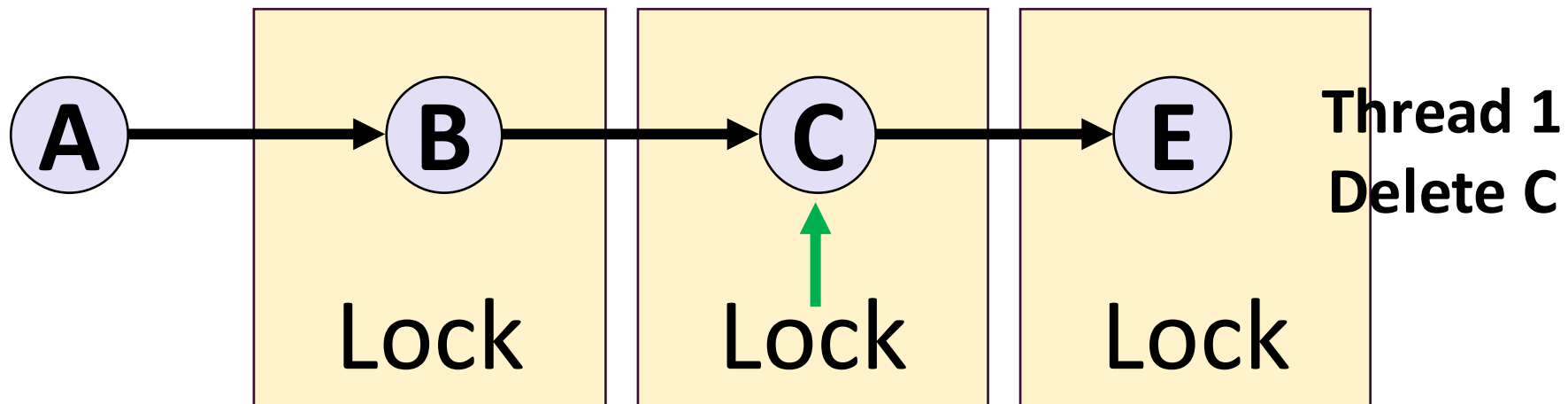
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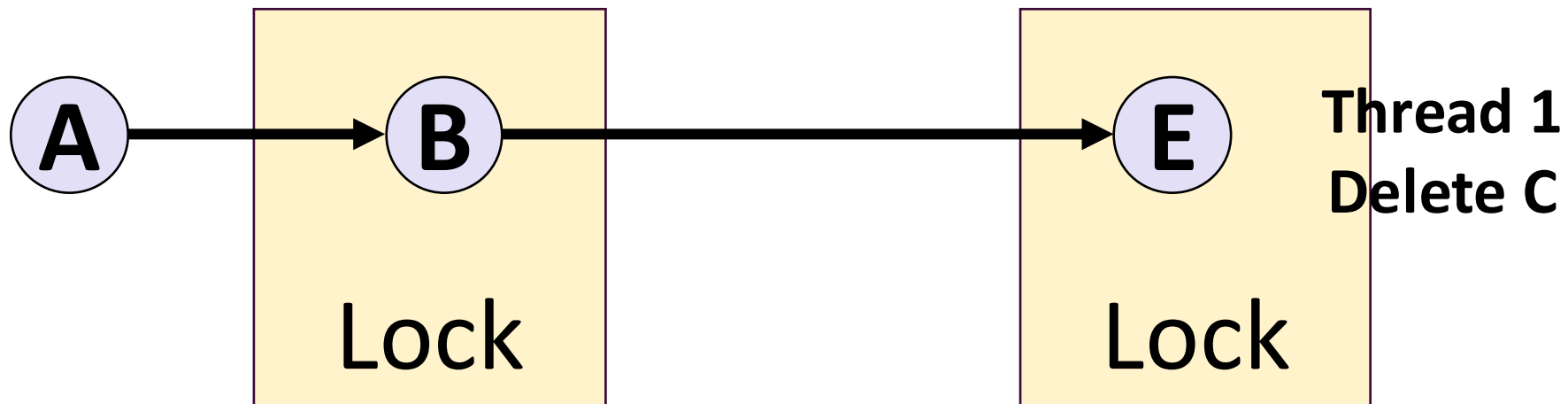
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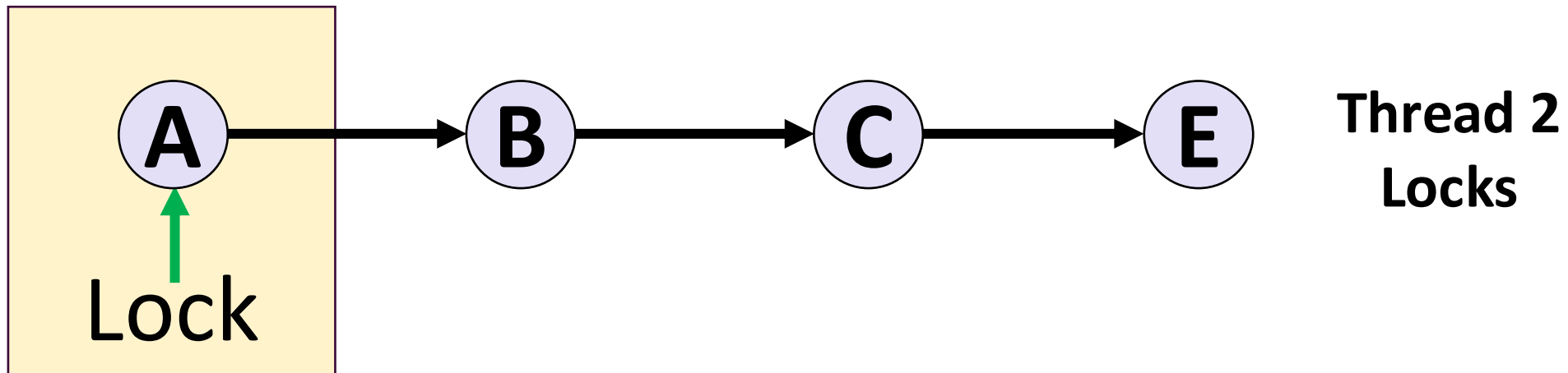
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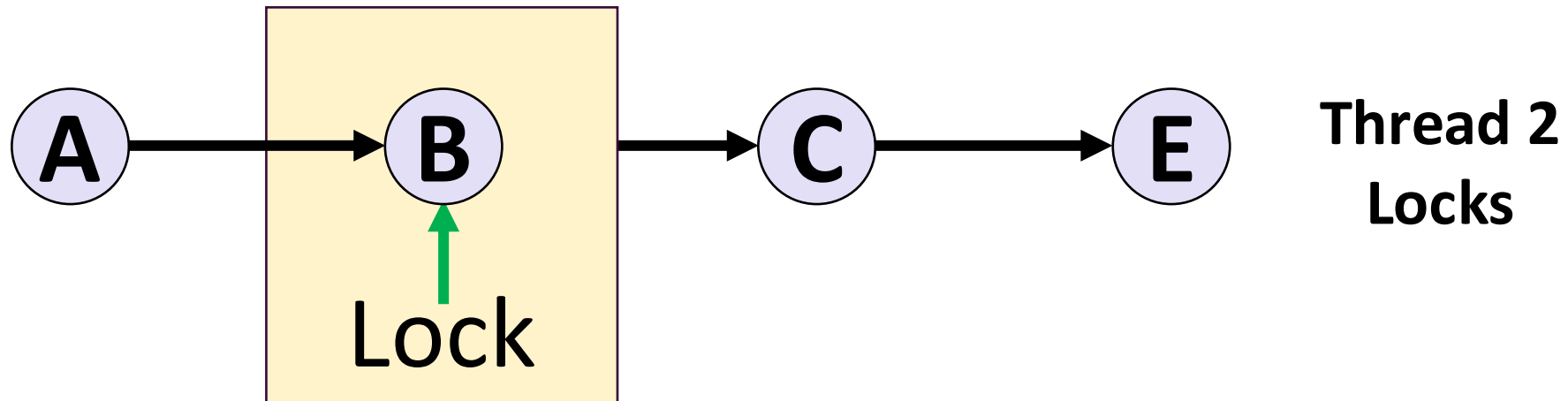
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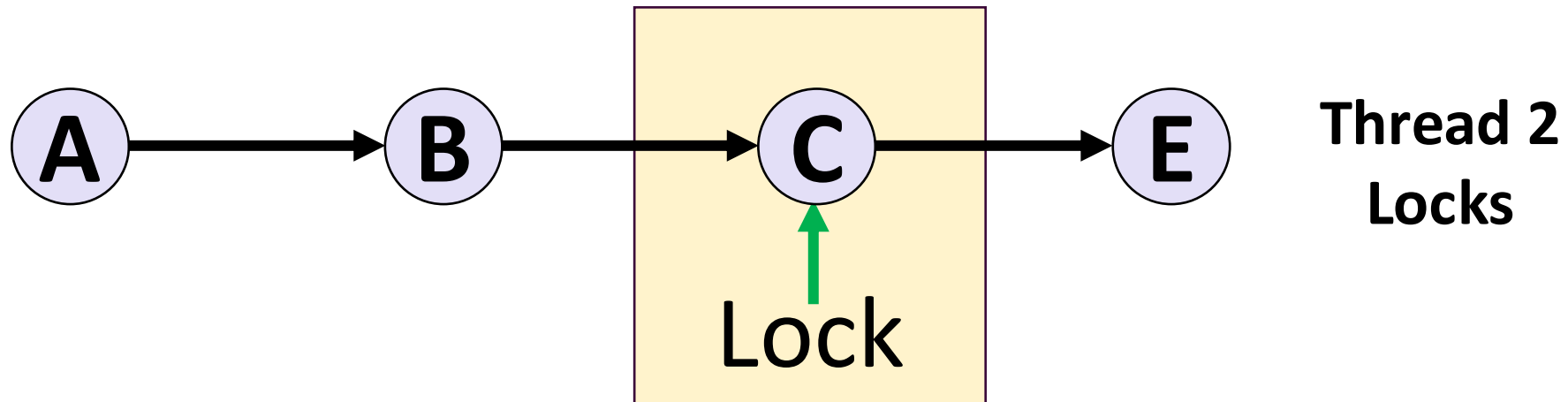
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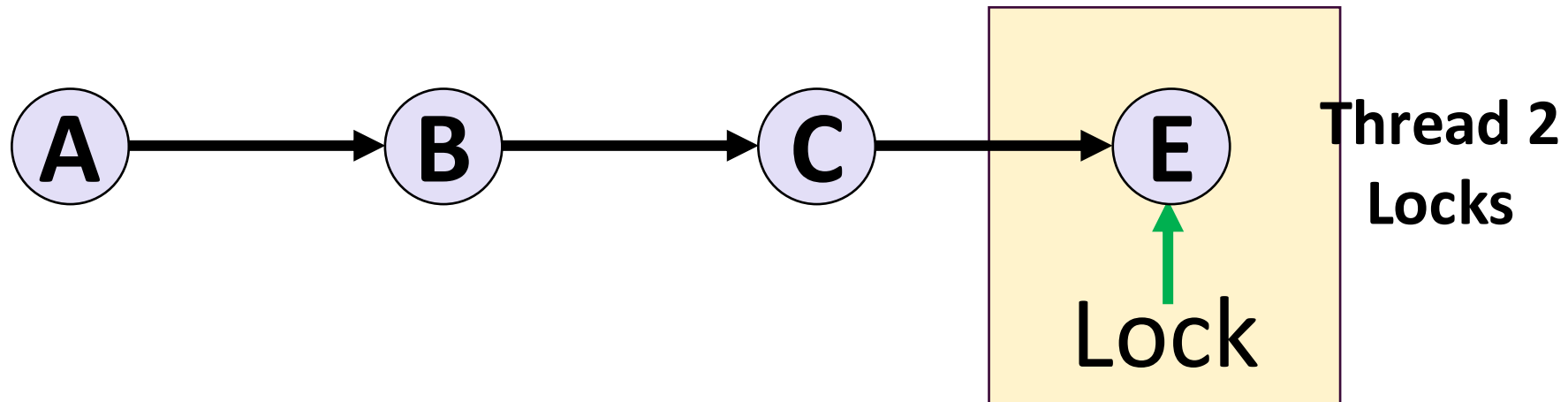
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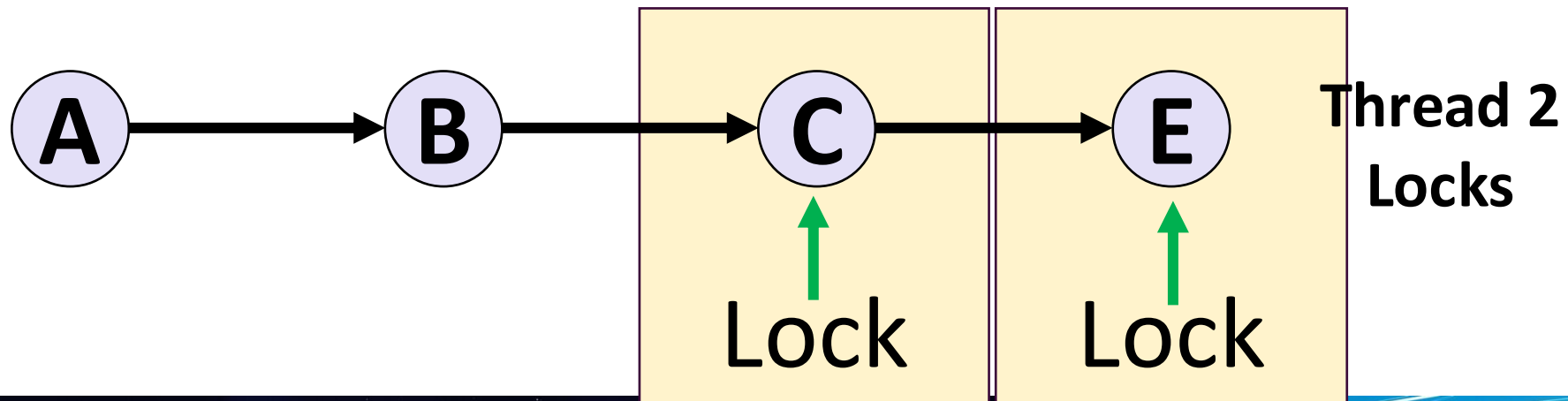
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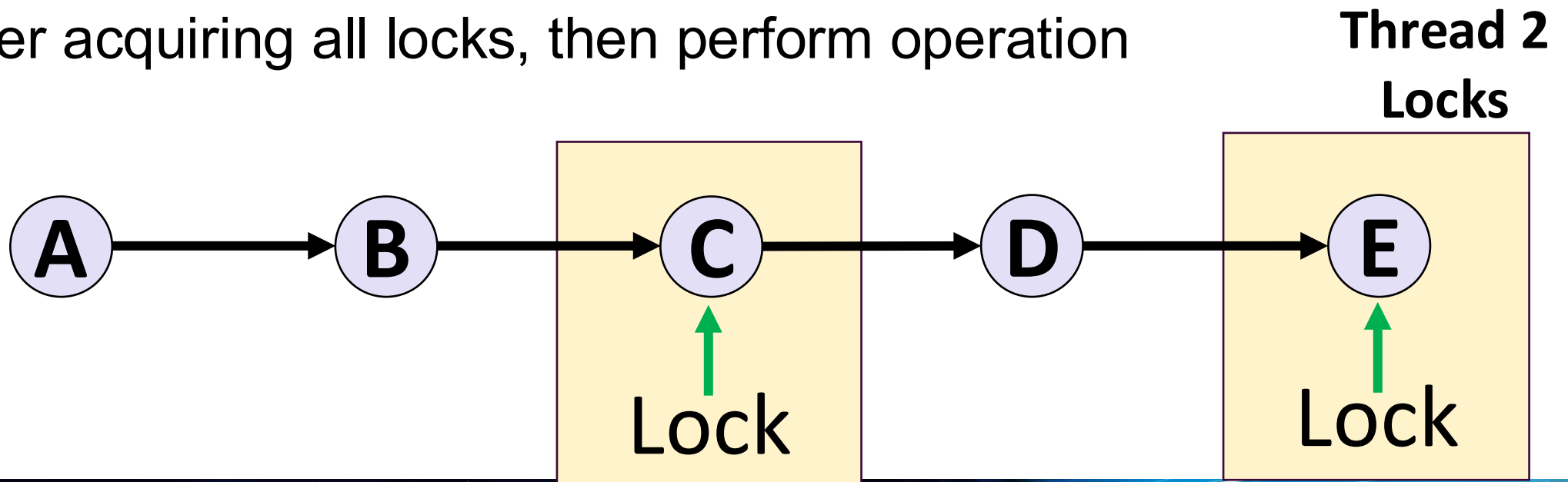
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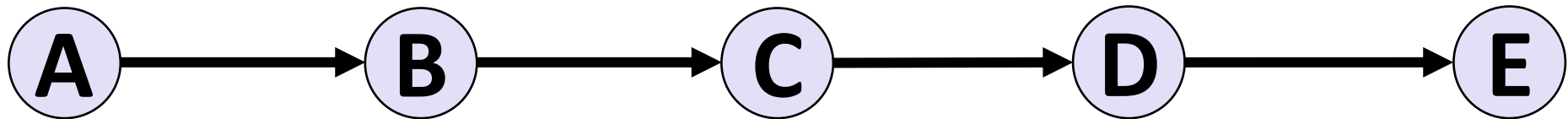
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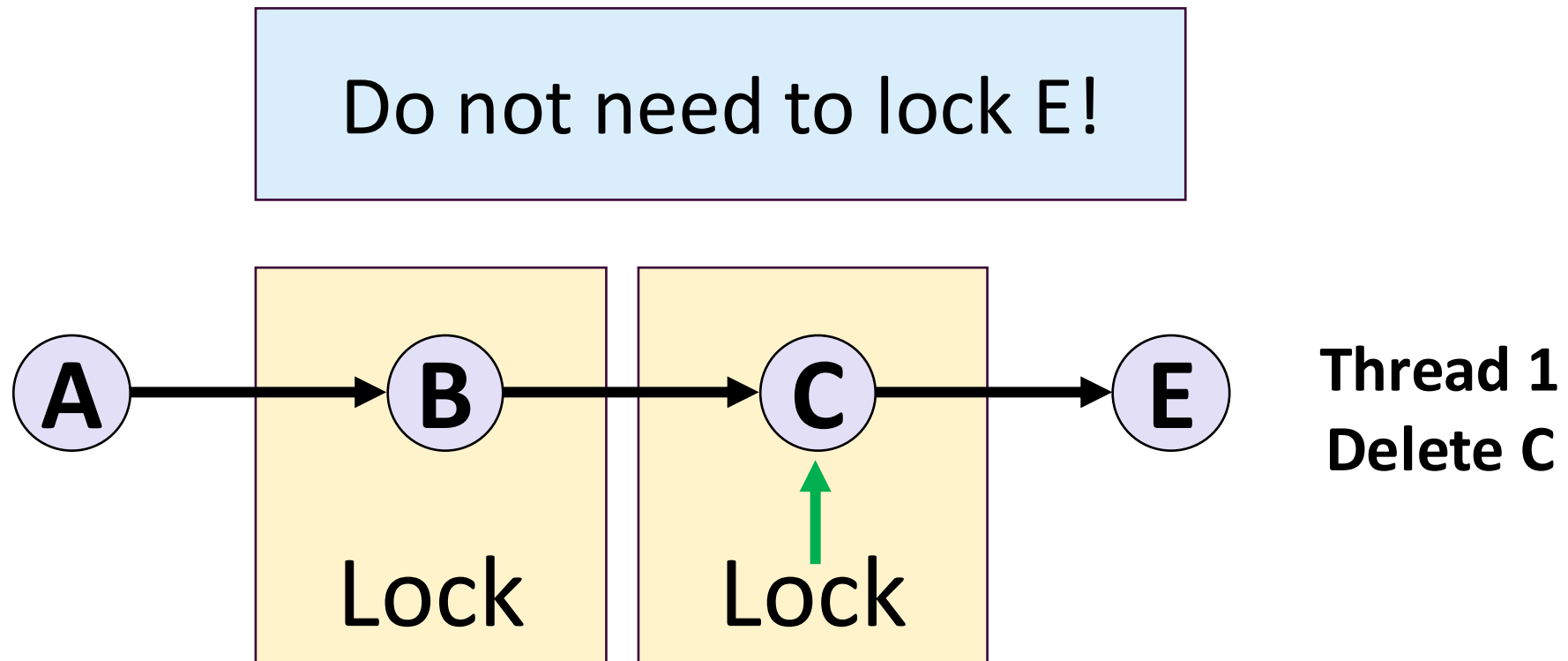
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Lock-Based Concurrent Linked List

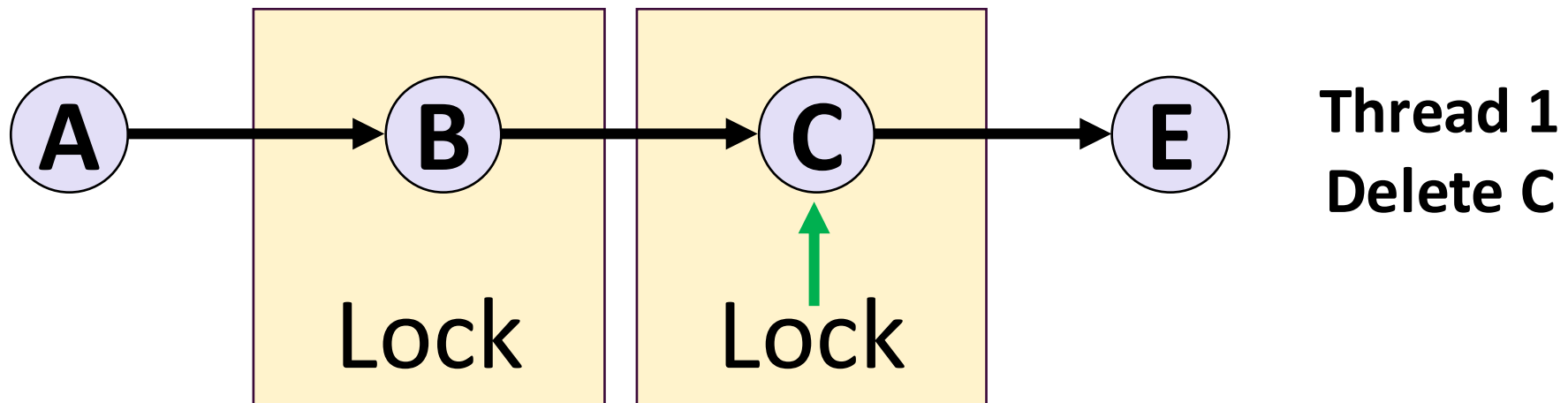
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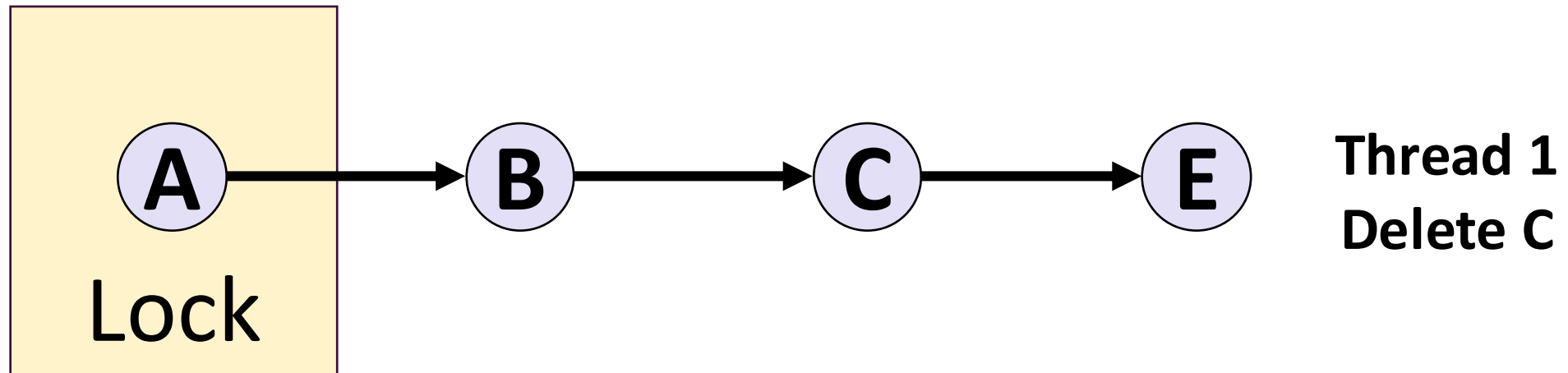
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Do not need to lock E!
We are only changing pointers of B



Lock-Based Concurrent Linked List

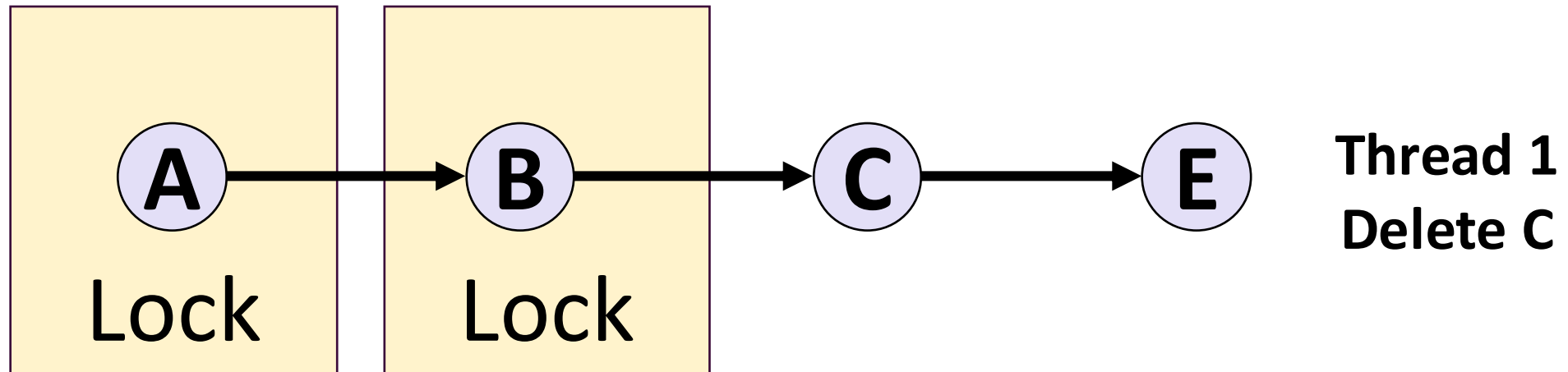
- **Revised procedure:**
 - Start at head and search for C



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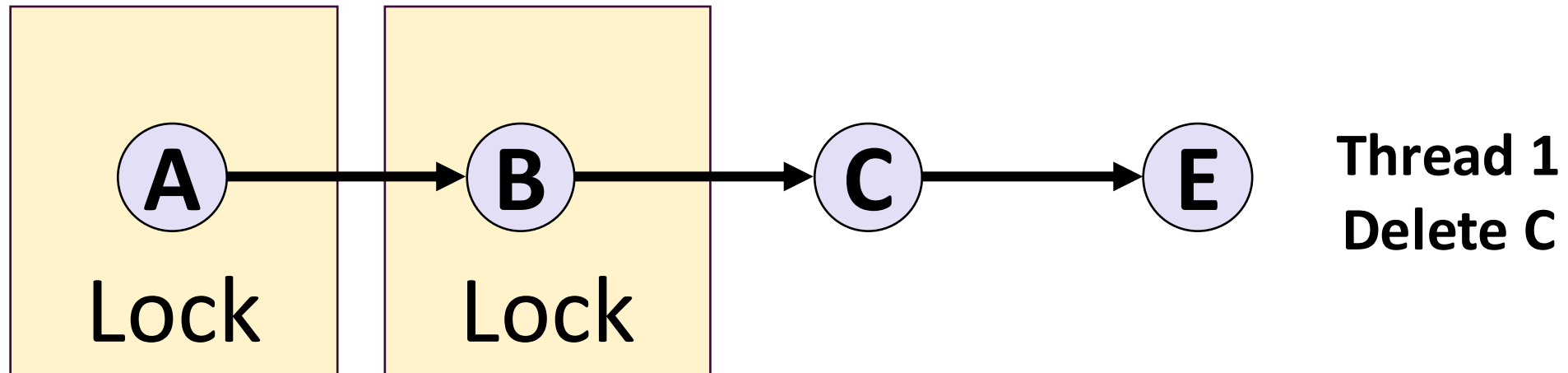
- Start at head and search for C
- Lock next searched node



Lock-Based Concurrent Linked List

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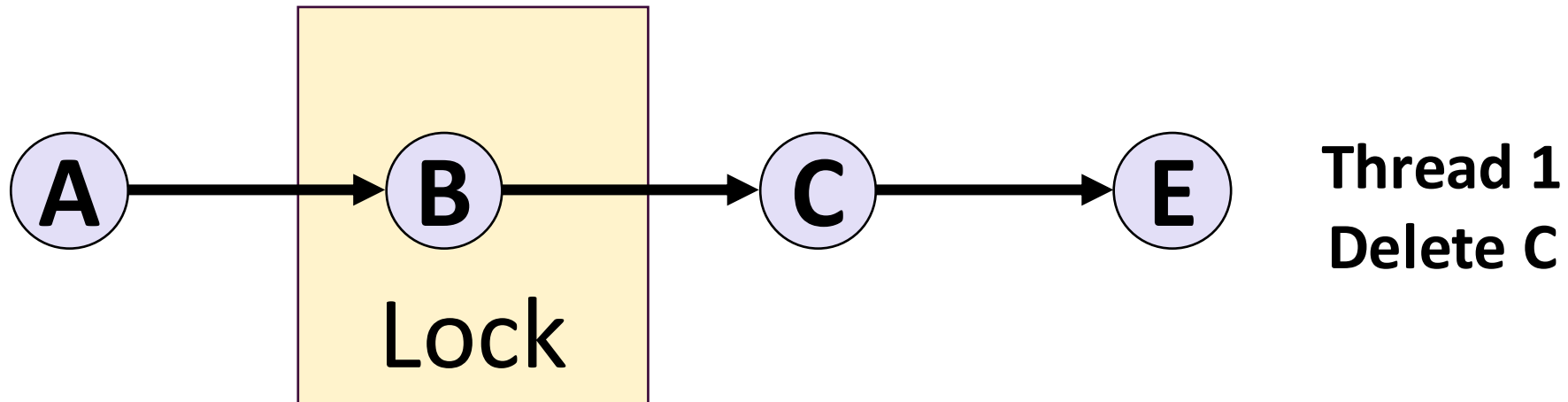
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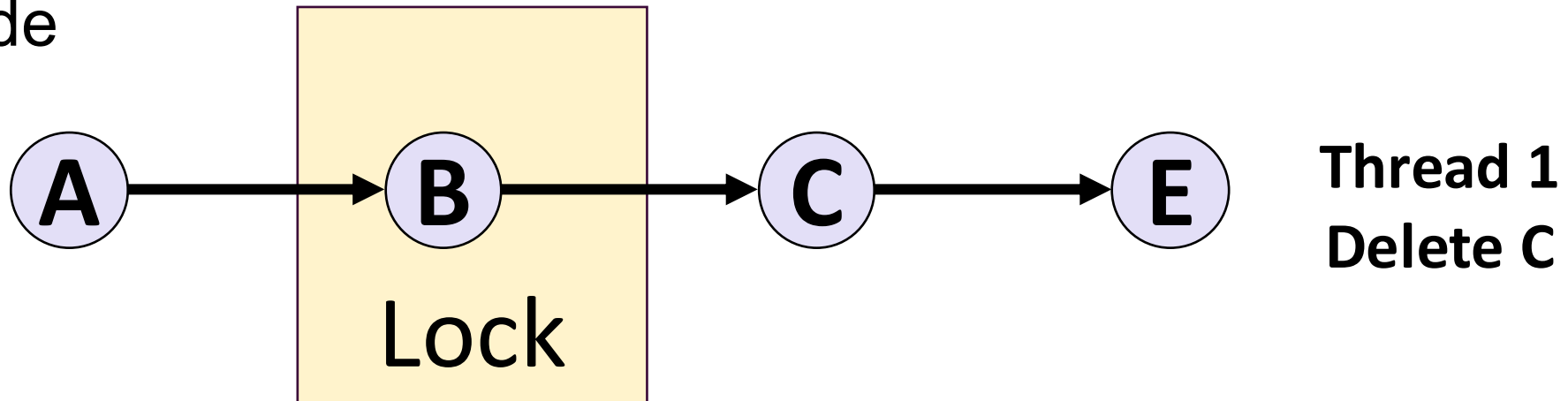
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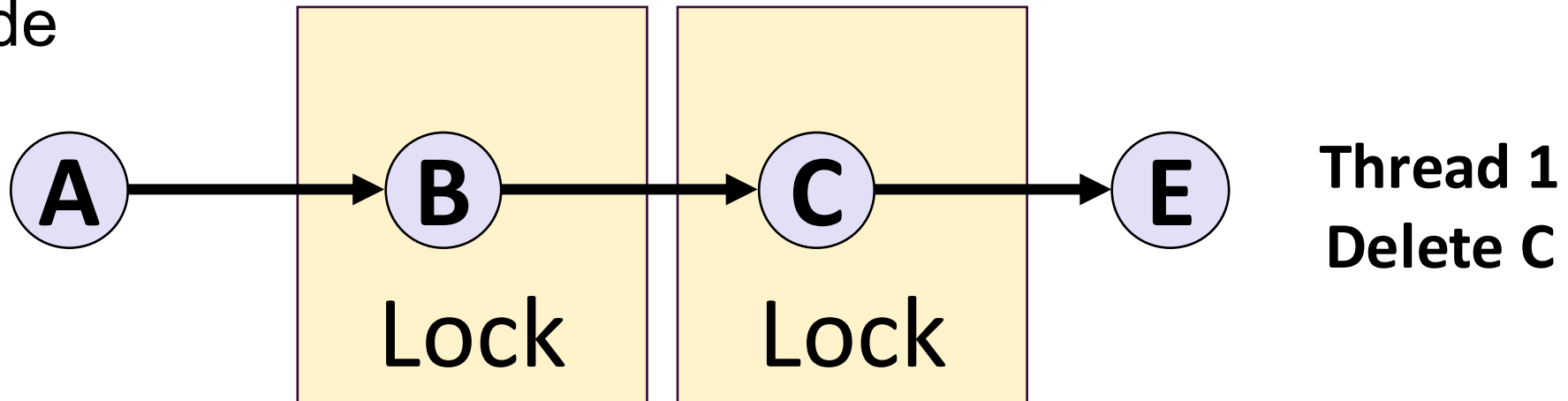
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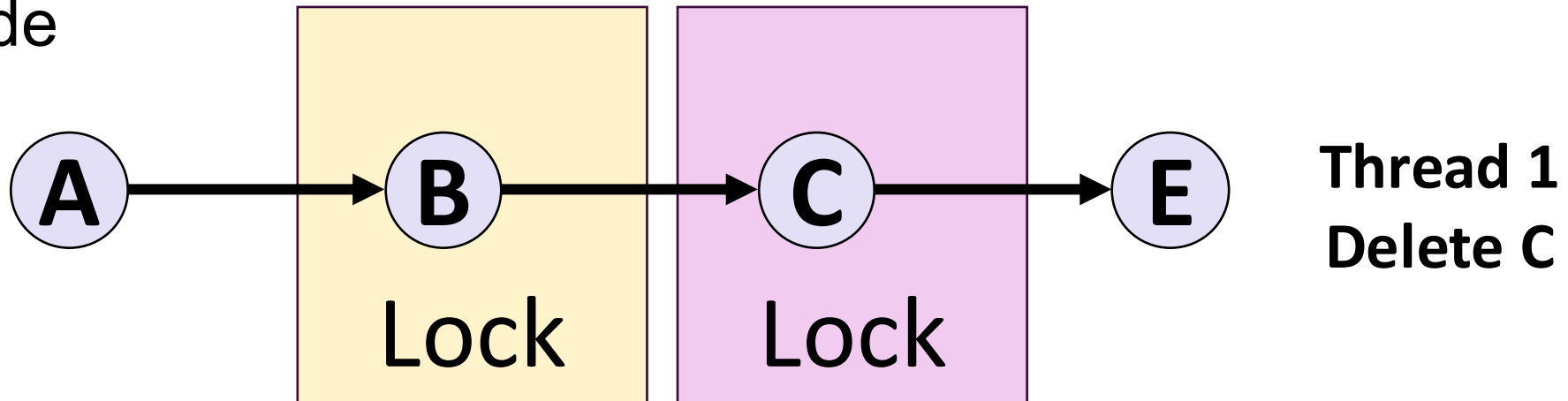
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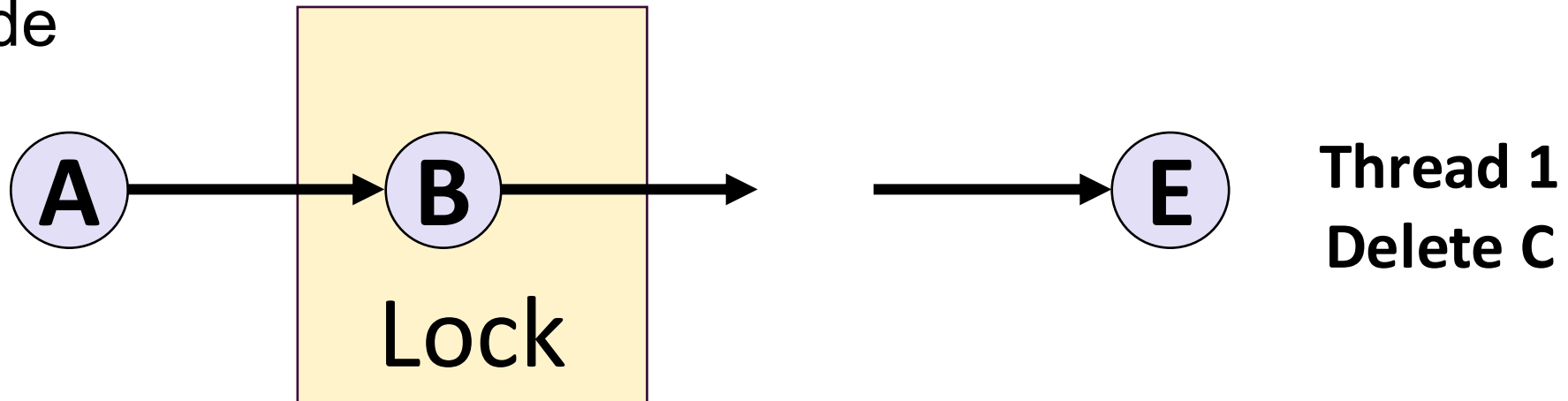
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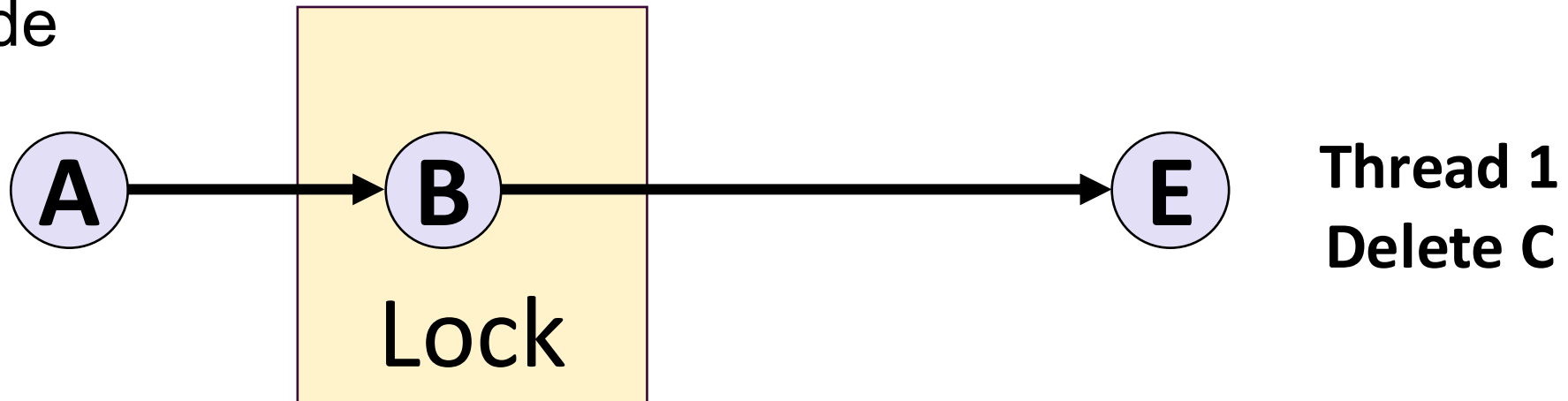
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Thread 1
Delete C

Lock-Based Concurrent Linked List

- **Revised procedure:**

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This is called “Hand-over-Hand” Locking!



**Thread 1
Delete C**

Lock-Based Concurrent Linked List

- **Thought exercise:**

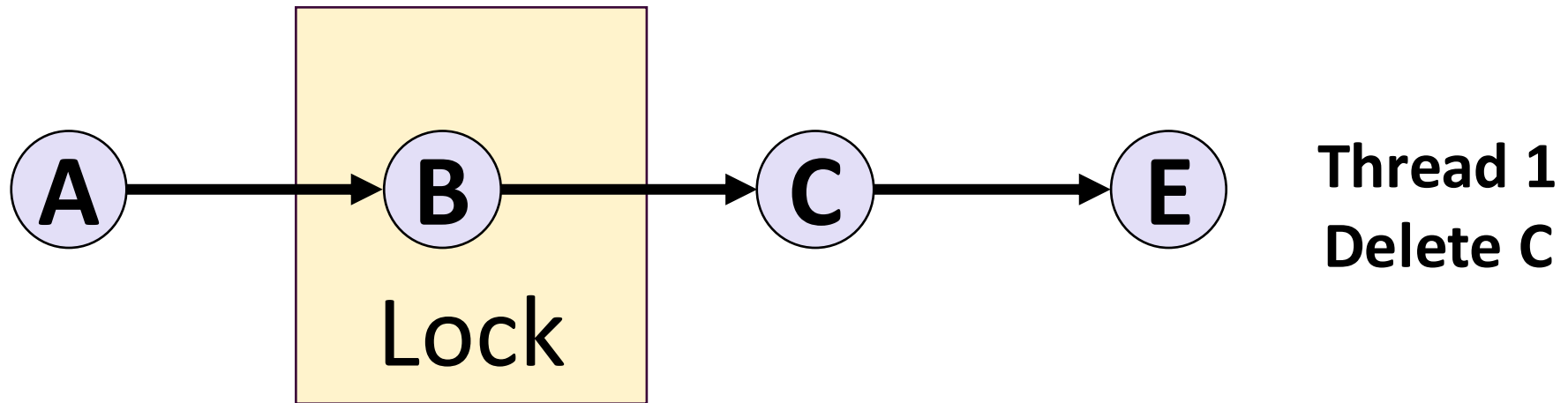
- What if we get rid of the lock on C when deleting C?
- **What could happen?**
- Is the lock on C necessary?



Thread 1
Delete C

Lock-Based Concurrent Linked List

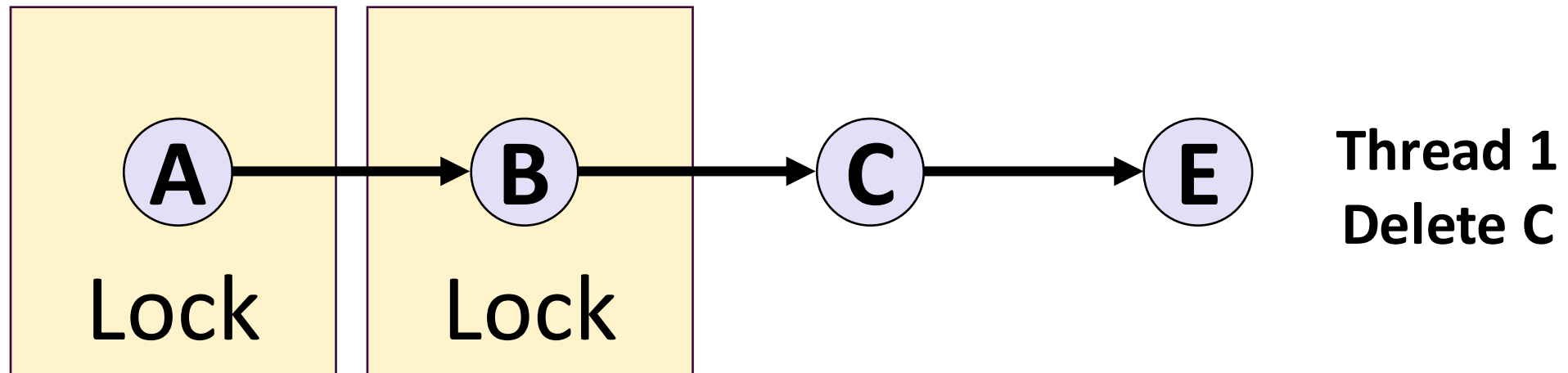
- There is a problem with not locking C!
- What is the problem?



Lock-Based Concurrent Linked List

- **Thought Experiment Procedure:**

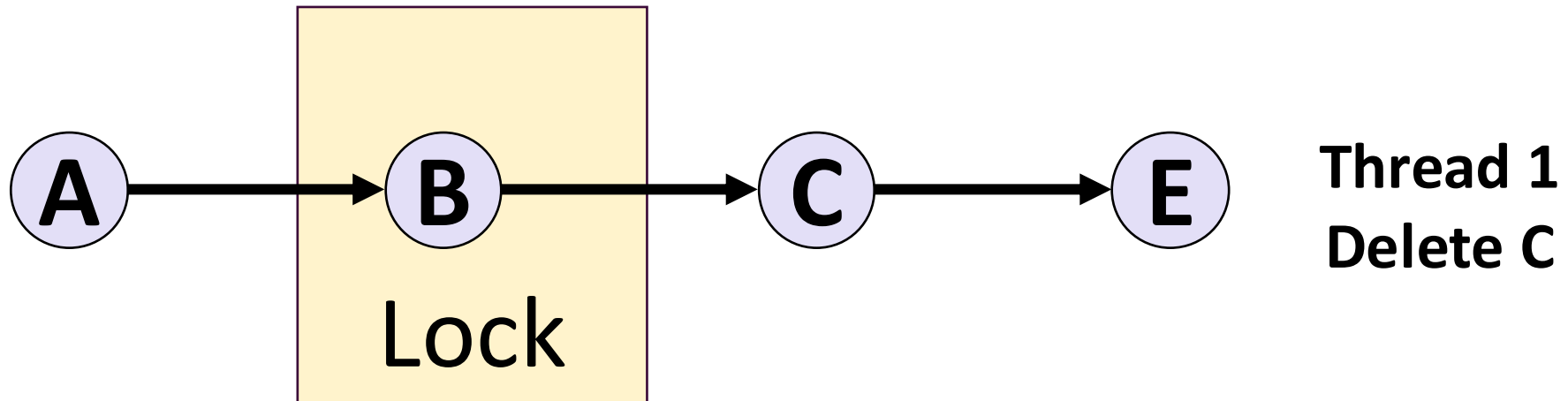
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Lock-Based Concurrent Linked List

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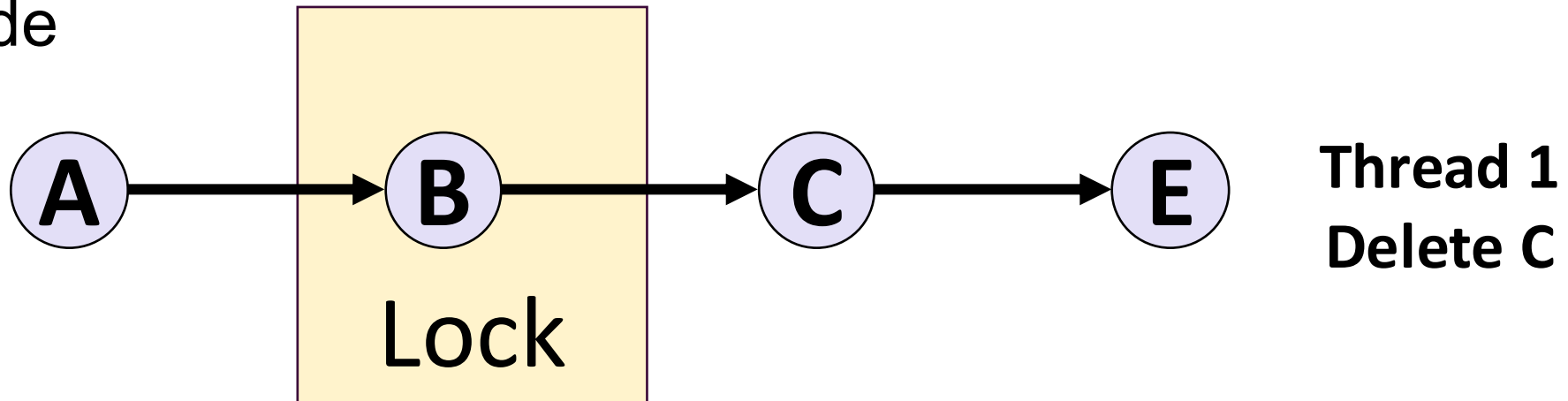
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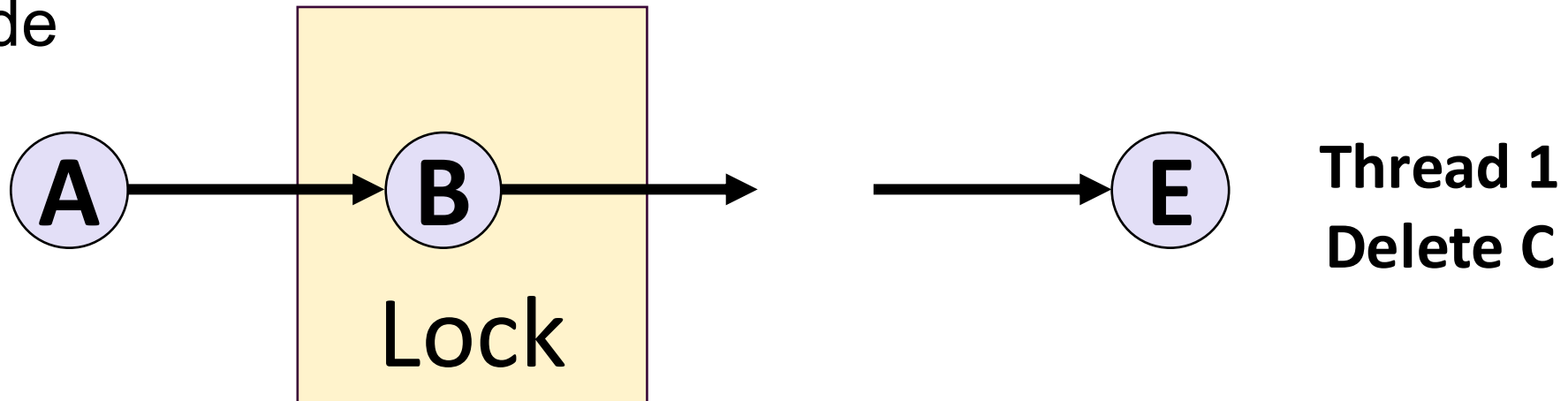
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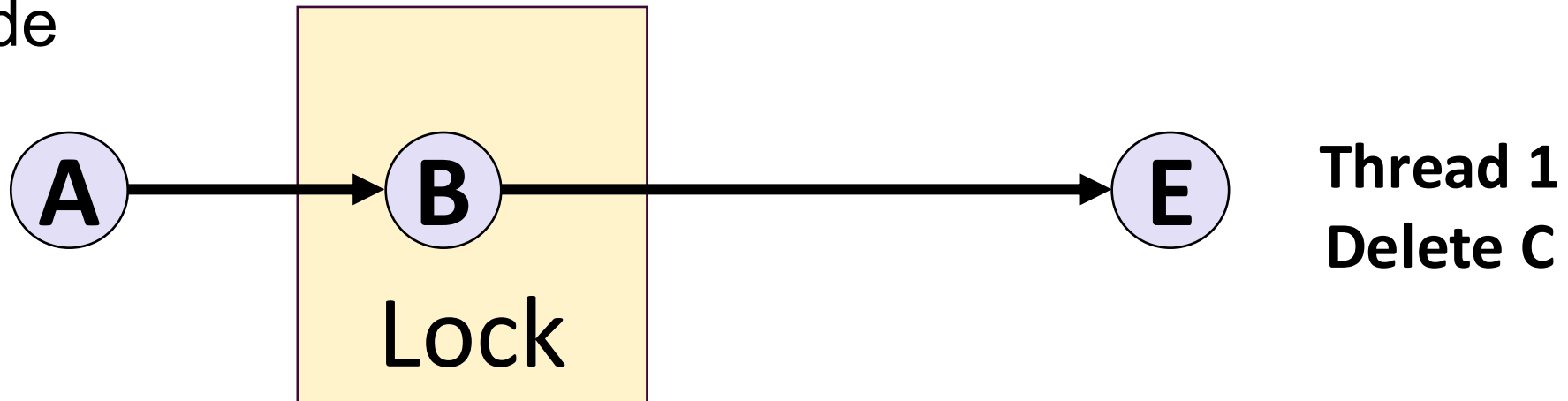
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Lock-Based Concurrent Linked List

- There is a problem with this thought experiment procedure!
- What is the problem?
- Suppose two threads, **one deletes B** and **one deletes C**
- We had discussed issues with concurrent algorithms before: deadlock, livelock, correctness issues...etc.

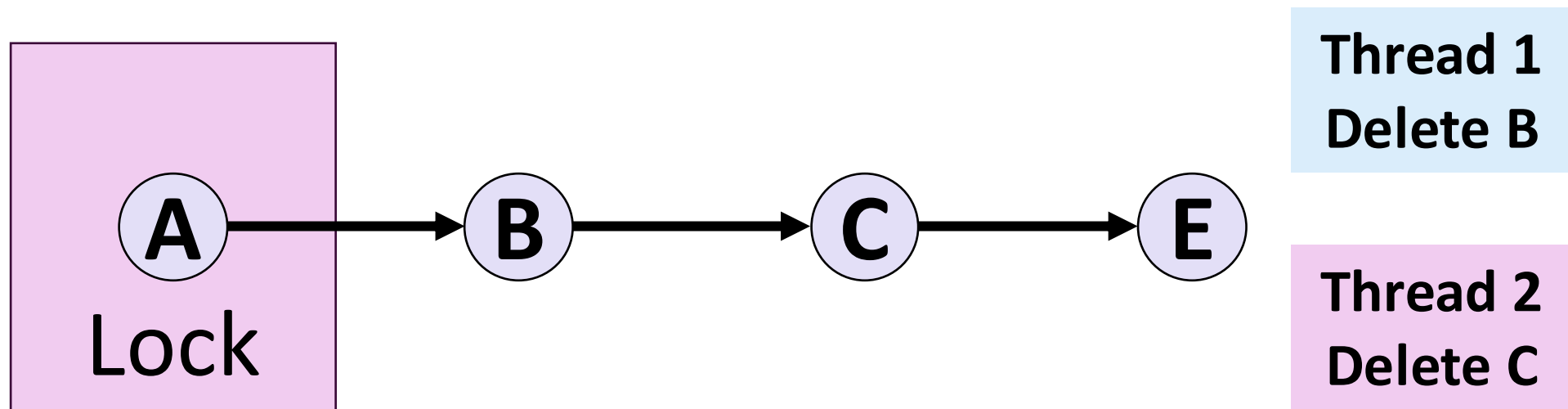


Thread 1
Delete B

Thread 2
Delete C

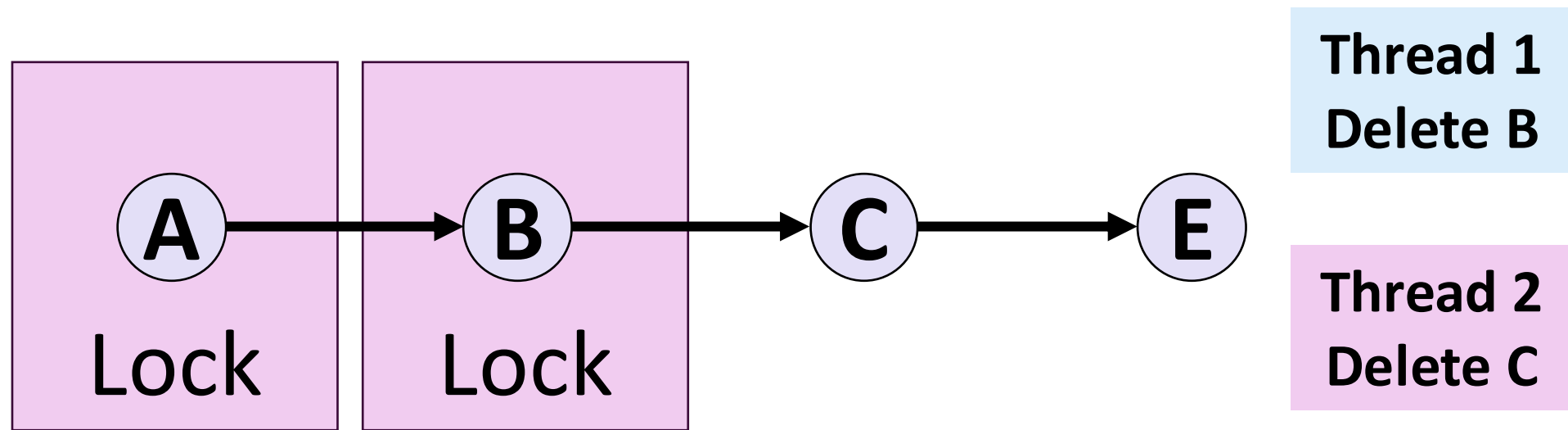
Lock-Based Concurrent Linked List Issue

- There could be deadlock caused by an incorrect state!



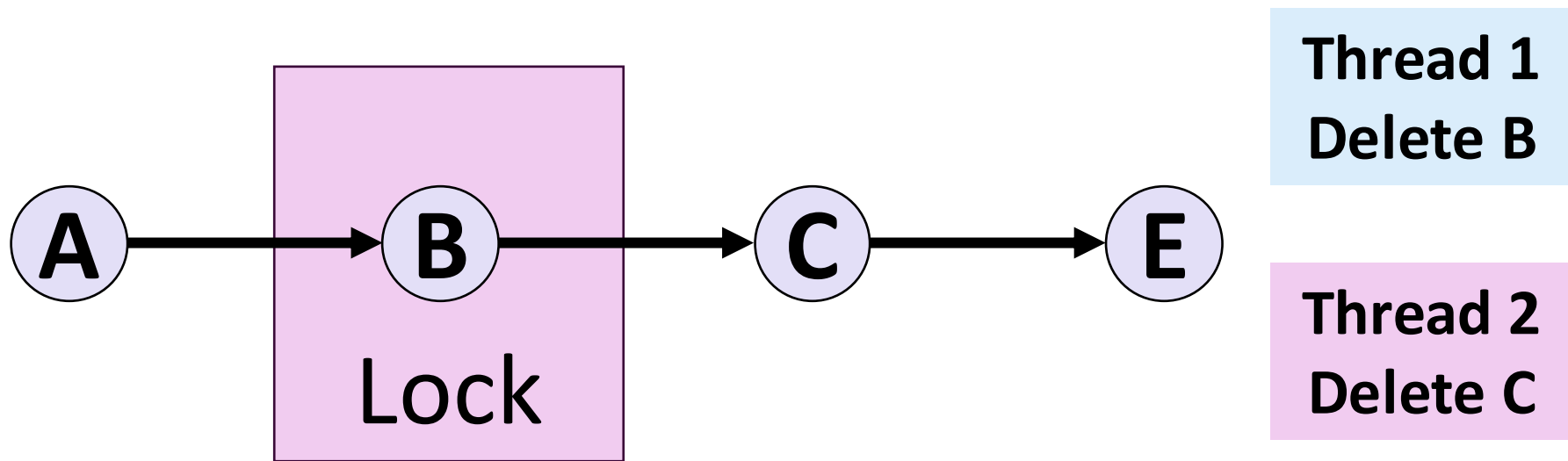
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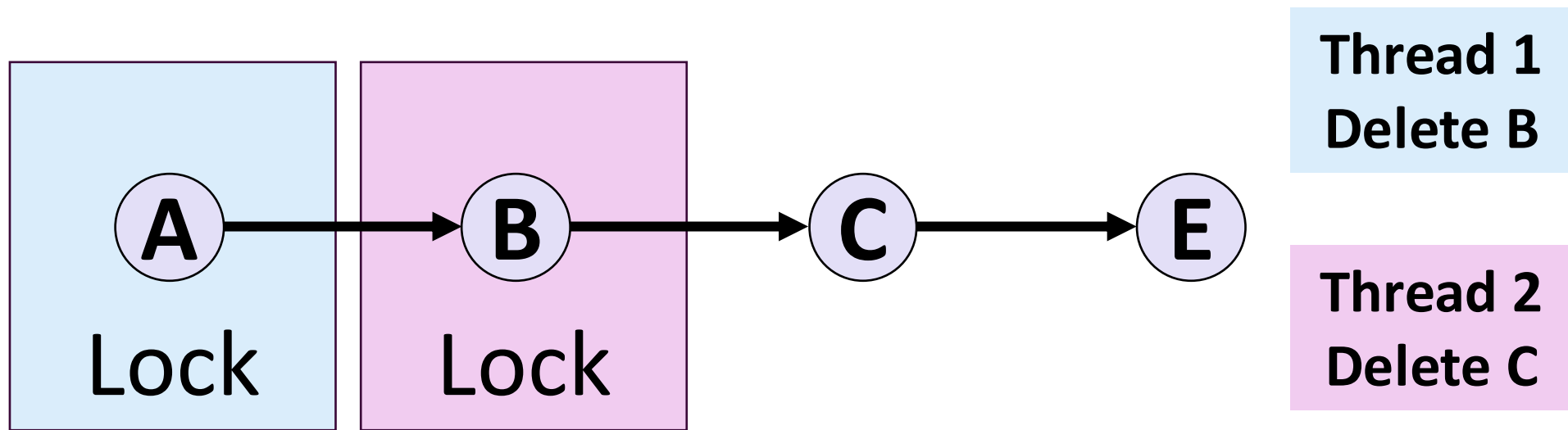
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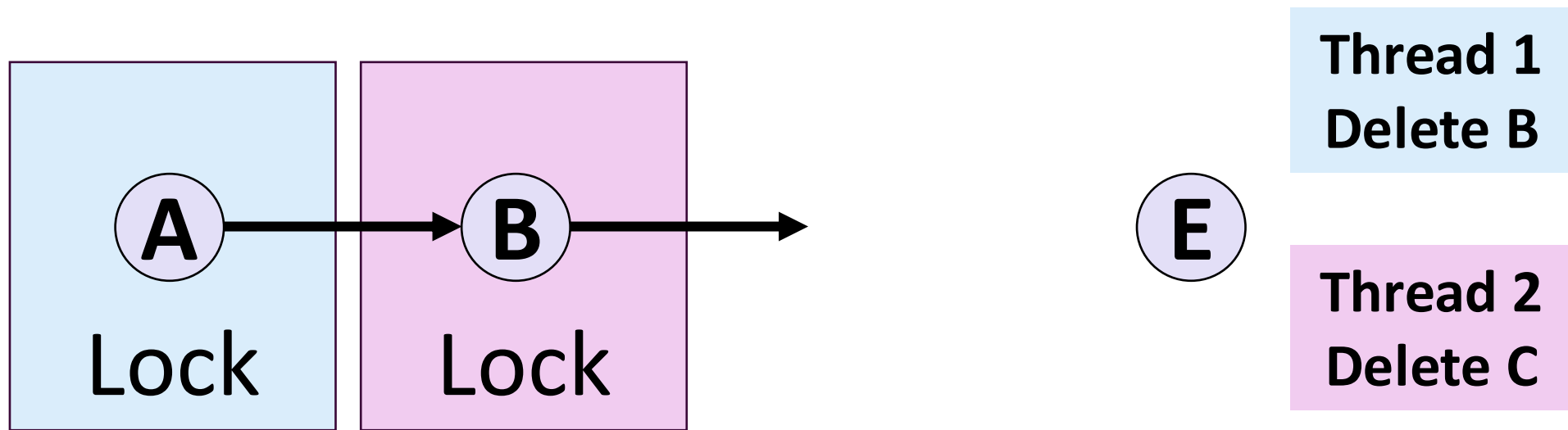
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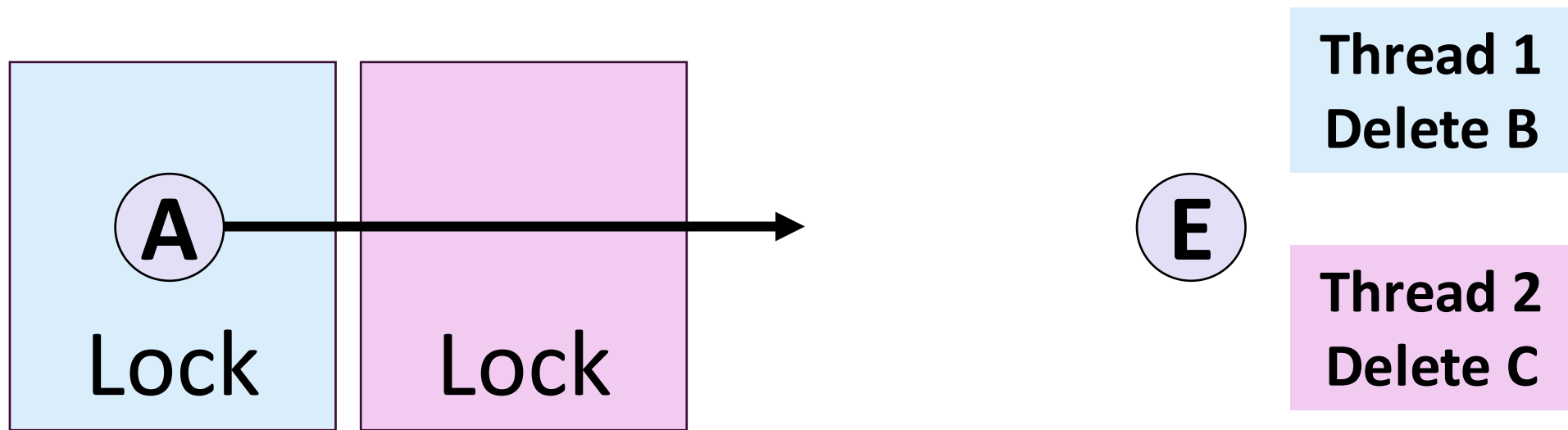
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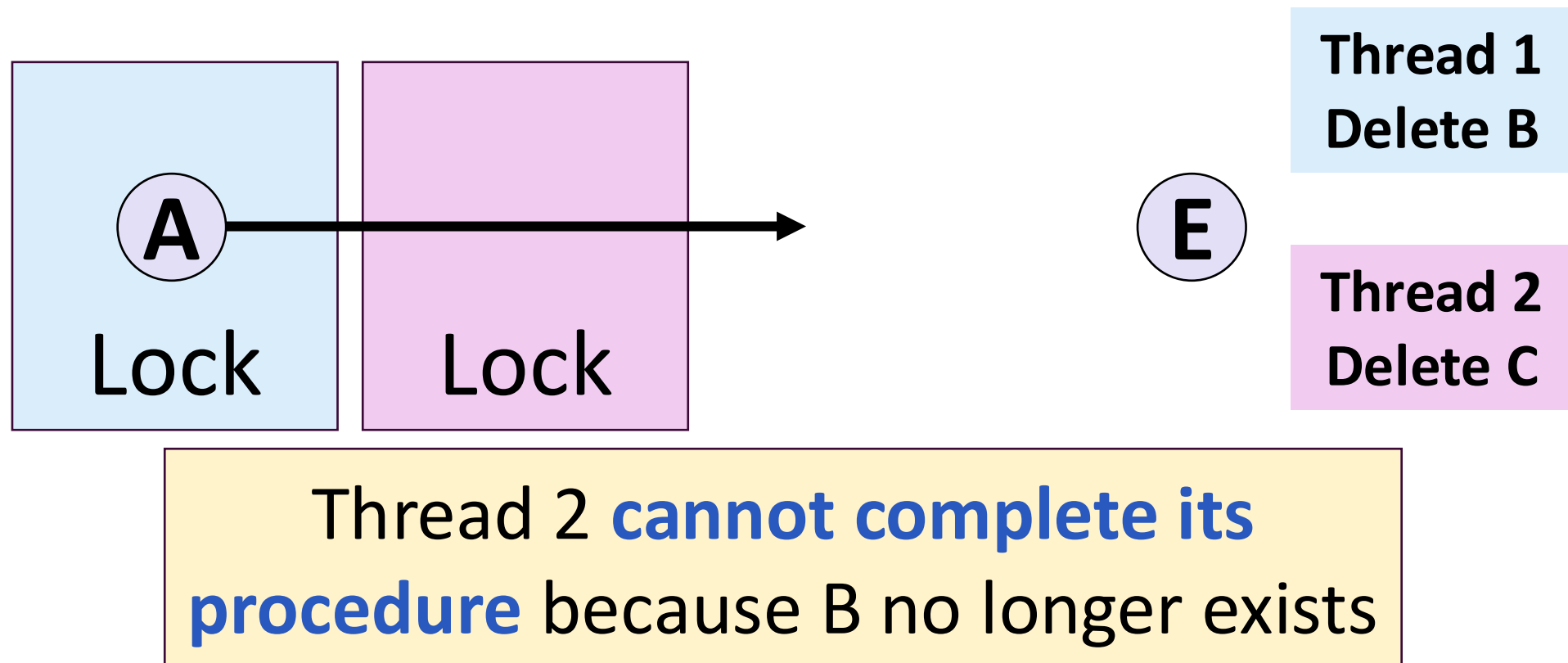
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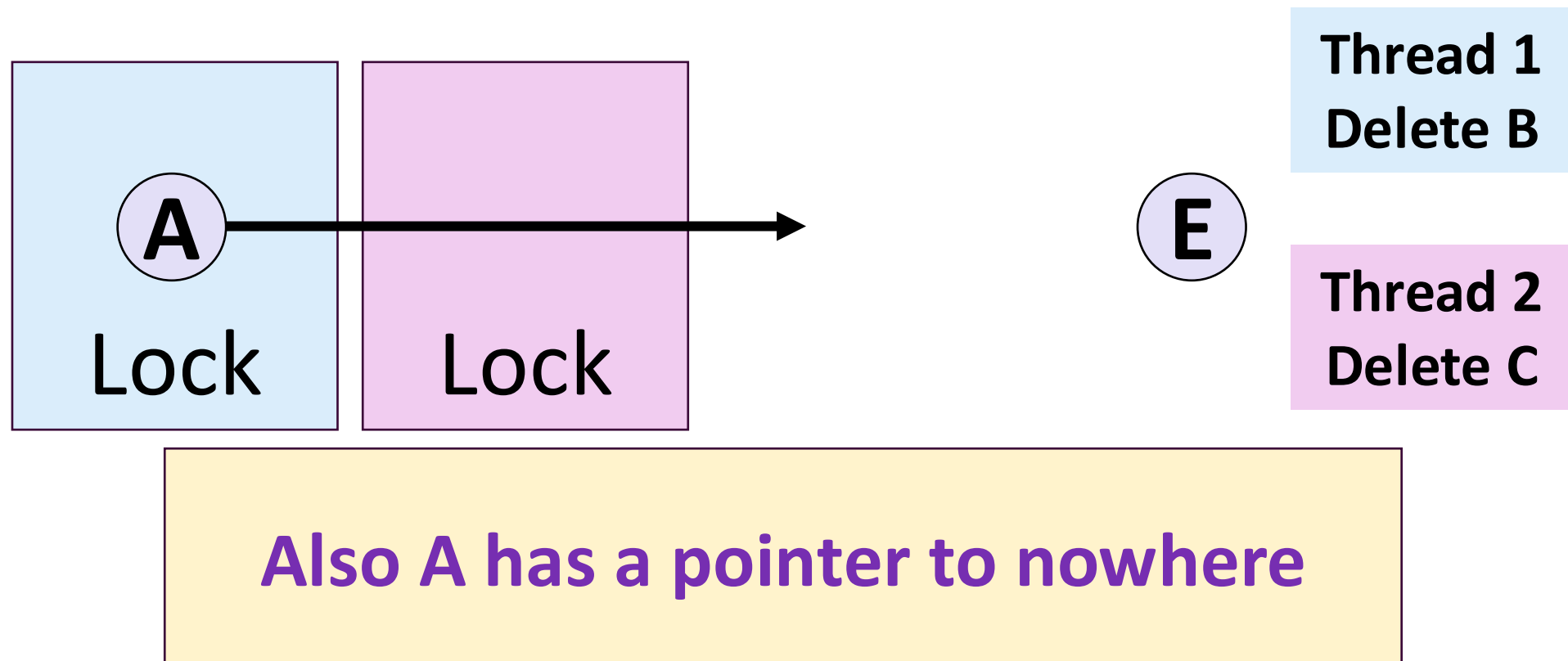
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Lock-Based Concurrent Linked List Issue

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Lock-Based Concurrent Linked List

- Despite being correct, this is not an ideal procedure still
- **What could cause suboptimal performance?**

Lock-Based Concurrent Linked List

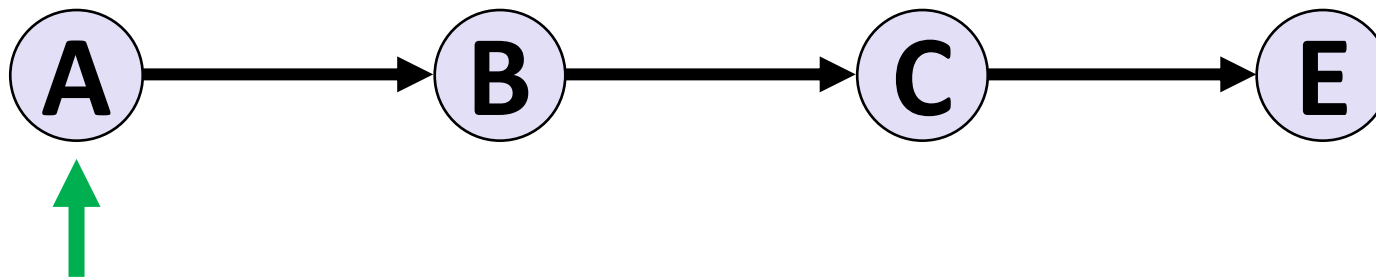
- Despite being correct, this is not an ideal procedure still
- Long chain of locking and unlocking
- Contention could cause this procedure to be very inefficient
- Serialized access to the procedure; deletions and insertions occur sequentially
- Increased latency from threads holding locks
 - A thread could take a long time to finish its operation; block other threads

Lock-Based Concurrent Linked List

- Currently, we do multiple locks:
 - Traversal requires a lock
 - Insertion requires three locks
 - Deletion requires two locks
- We can try to reduce the number of locks

Optimistic Locking

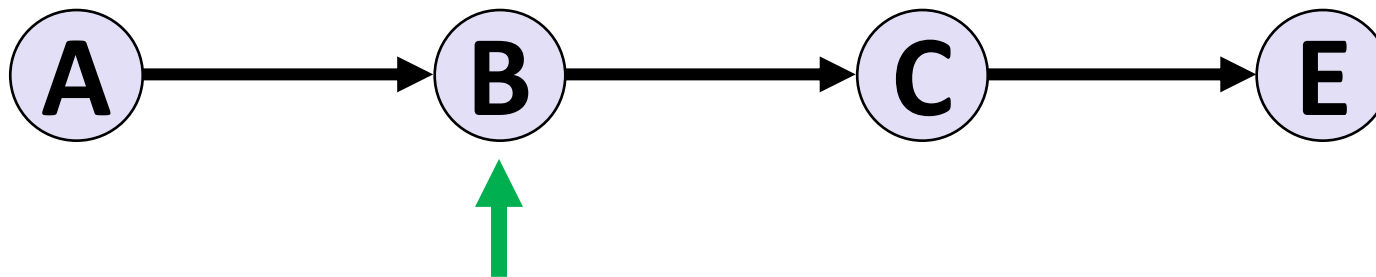
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- **Only lock when you find the correct nodes**
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Thread 1
Delete C

Optimistic Locking

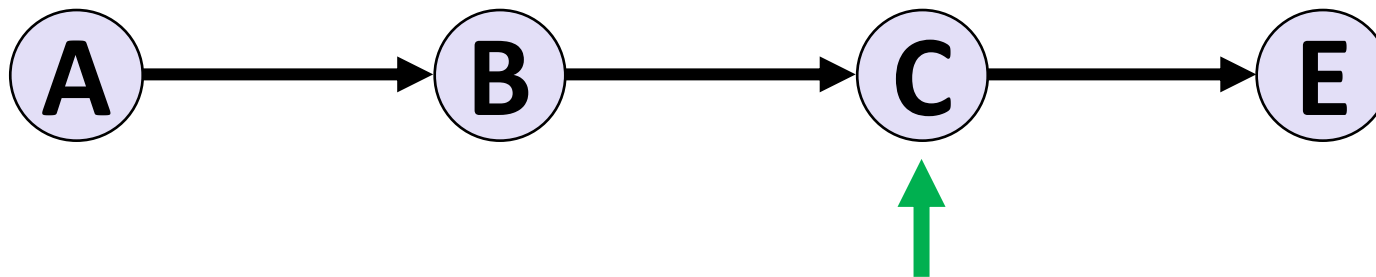
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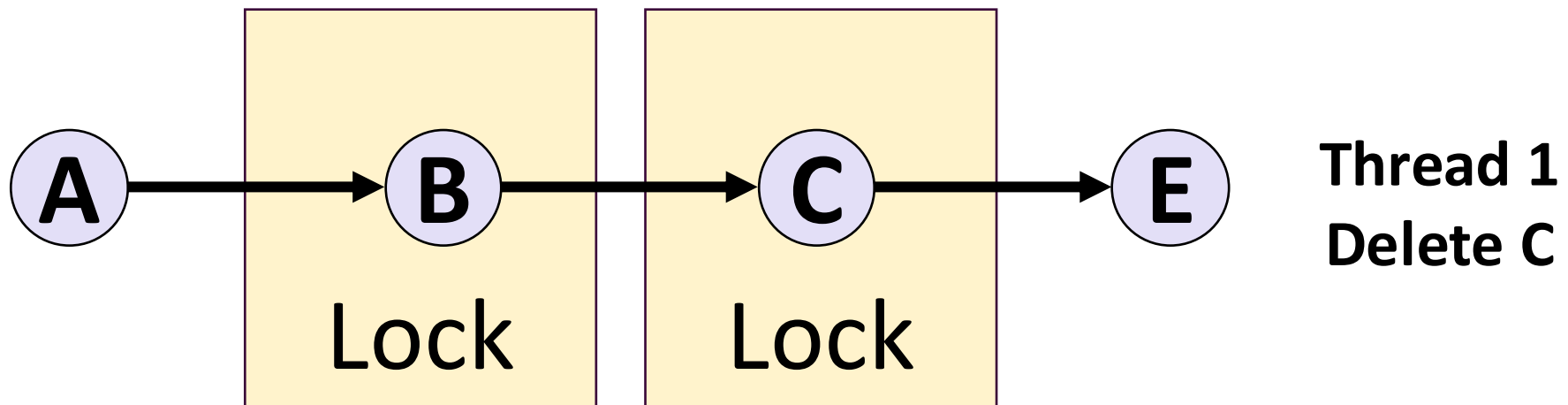
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Thread 1
Delete C

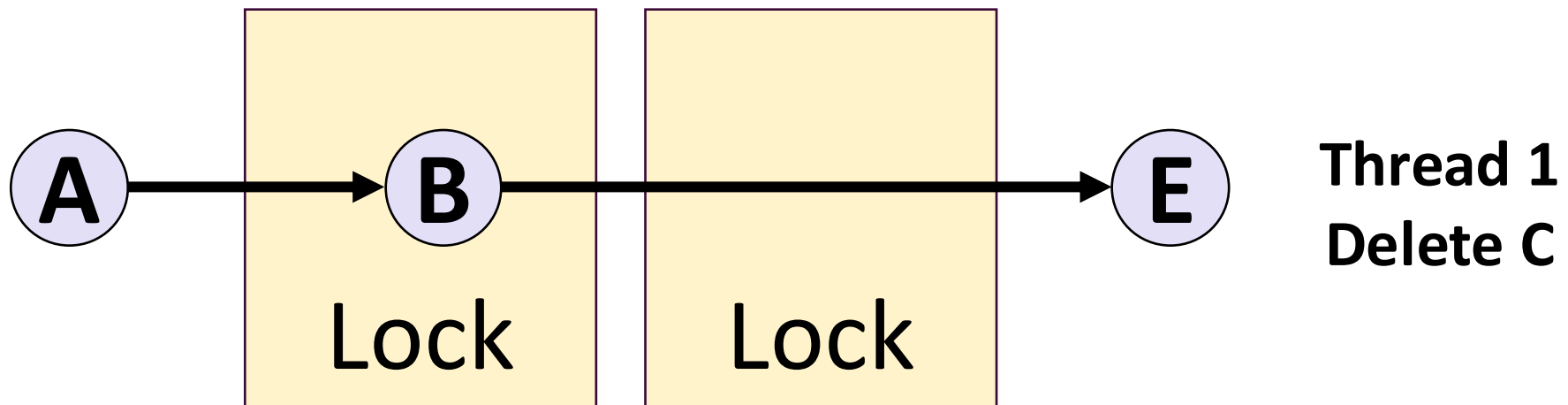
Optimistic Locking

- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
- Finally, release the locks



Optimistic Locking

- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
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Optimistic Locking

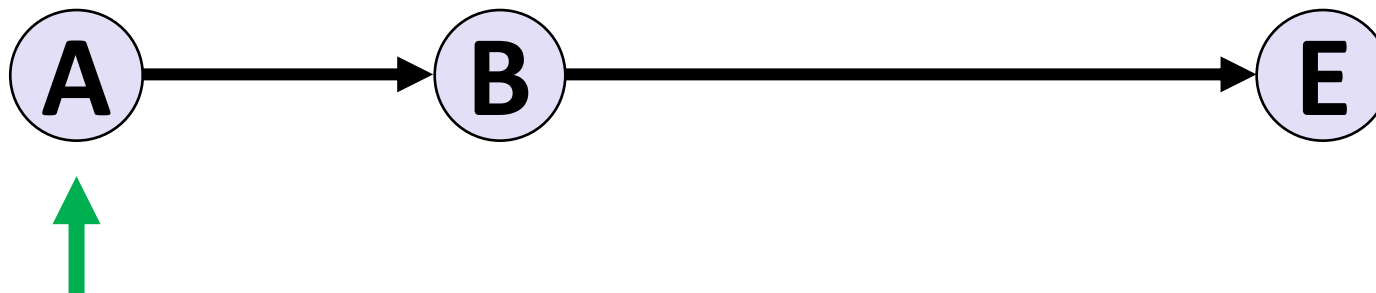
- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
- Finally, release the locks



Thread 1
Delete C

Optimistic Locking

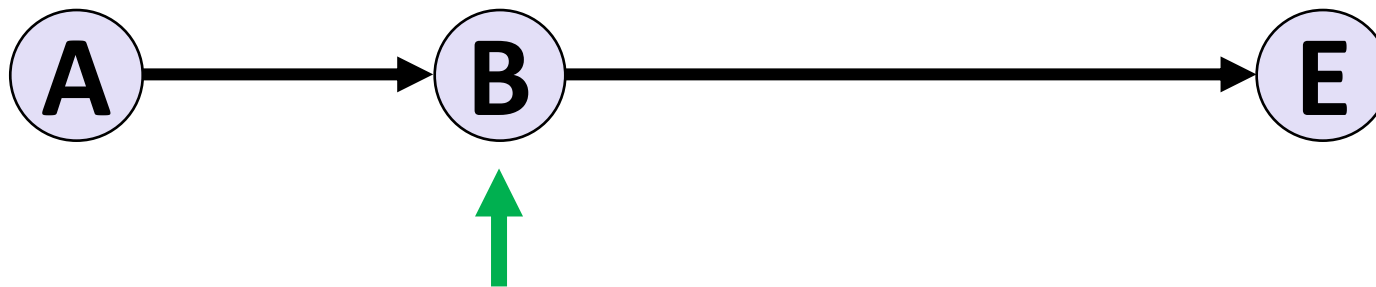
- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
- Finally, release the locks



Thread 2
Insert D

Optimistic Locking

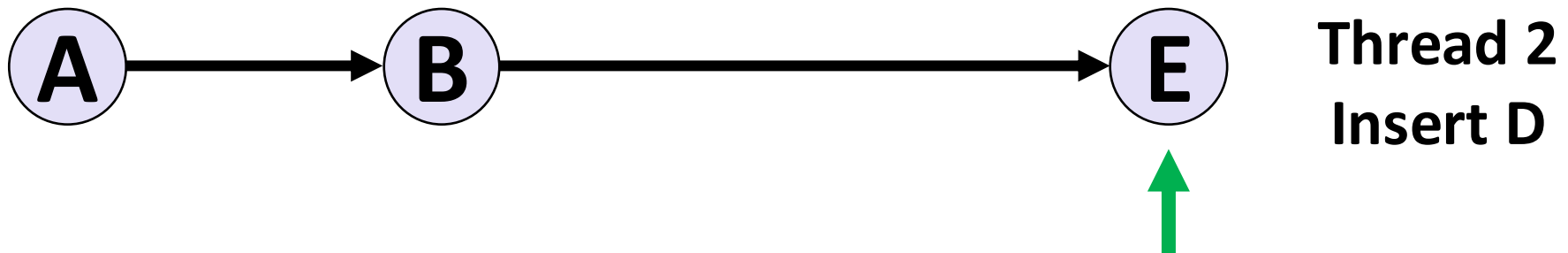
- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
- Finally, release the locks



Thread 2
Insert D

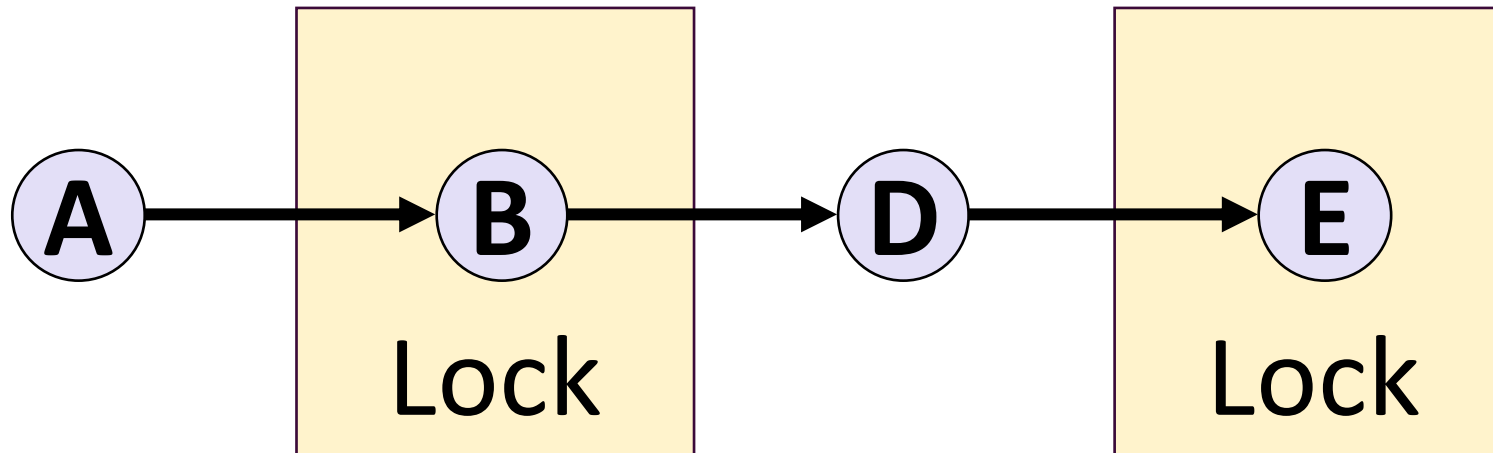
Optimistic Locking

- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
- Finally, release the locks



Optimistic Locking

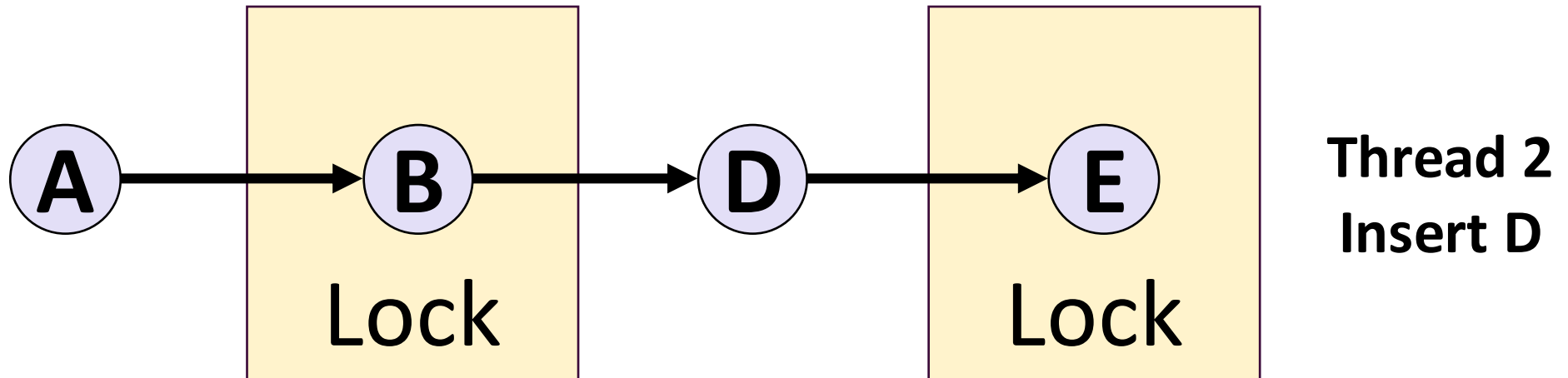
- Be optimistic and **traverse without locking**
- **Only lock when you find the correct nodes**
- Once you have the locks, then perform the update
- Finally, release the locks



Thread 2
Insert D

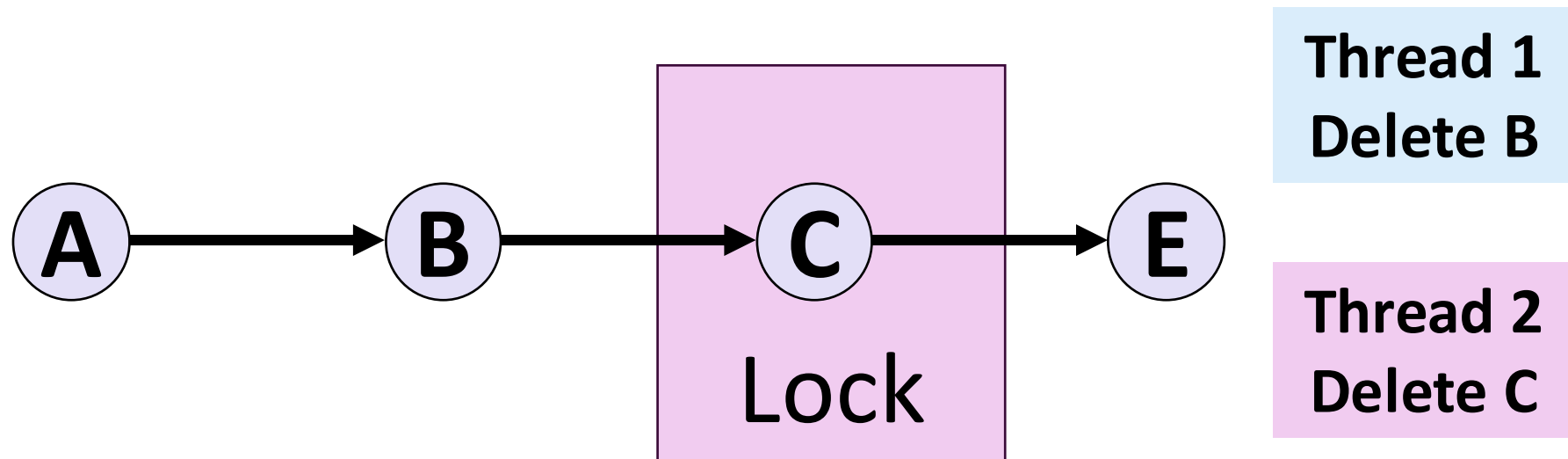
Optimistic Locking

- Can anything go wrong with this new procedure?



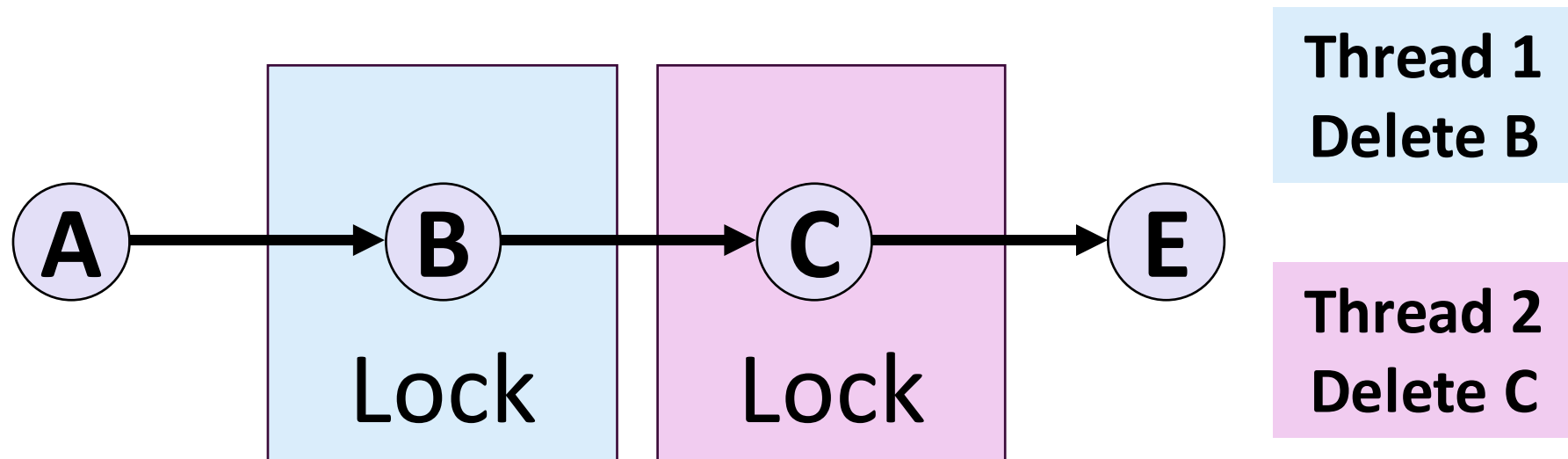
Optimistic Locking

- Can anything go wrong with this new procedure?
- First, need to acquire locks on the predecessor then successor in order to prevent deadlocks



Optimistic Locking

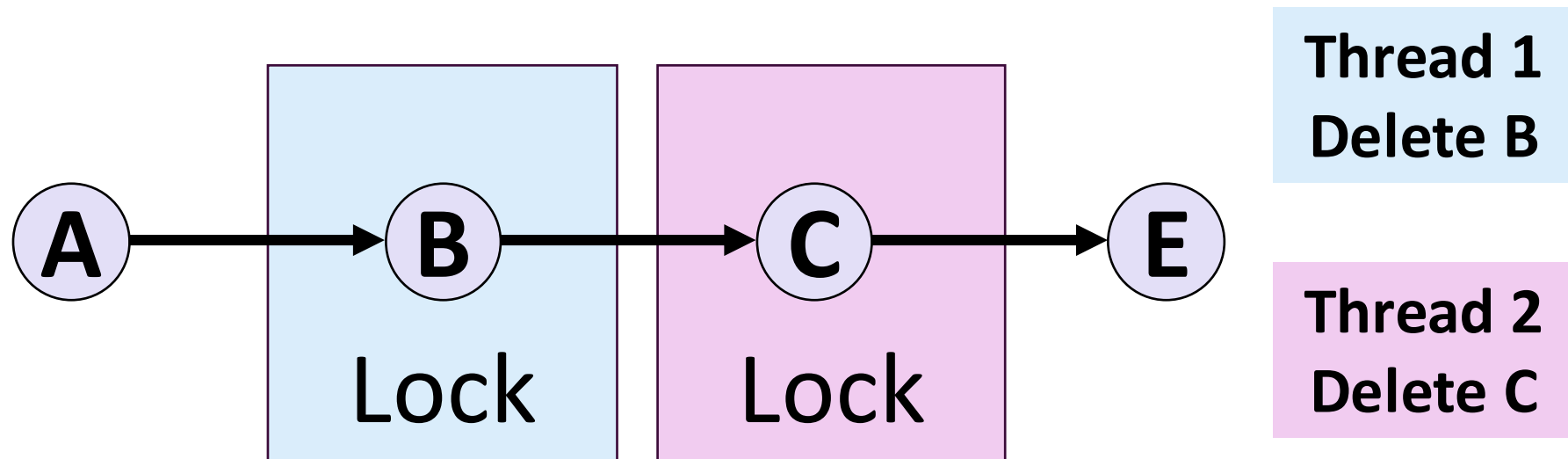
- Can anything go wrong with this new procedure?
- First, need to acquire locks on the predecessor then successor in order to prevent deadlocks



Optimistic Locking

- Can anything go wrong with this new proc
- First, need to acquire locks on the predecessor successor in order to prevent deadlocks

Deadlock! When Thread 1 attempts to lock C and Thread 2 attempts to lock B



Optimistic Locking

- Can anything go wrong with this new procedure?
- First, need to acquire locks on the predecessor then successor in order to prevent deadlocks
- Anything else?
 - What if you have a simultaneous insertion and deletion?

Optimistic Locking

- Anything else?

- What if you have a simultaneous insertion and deletion?



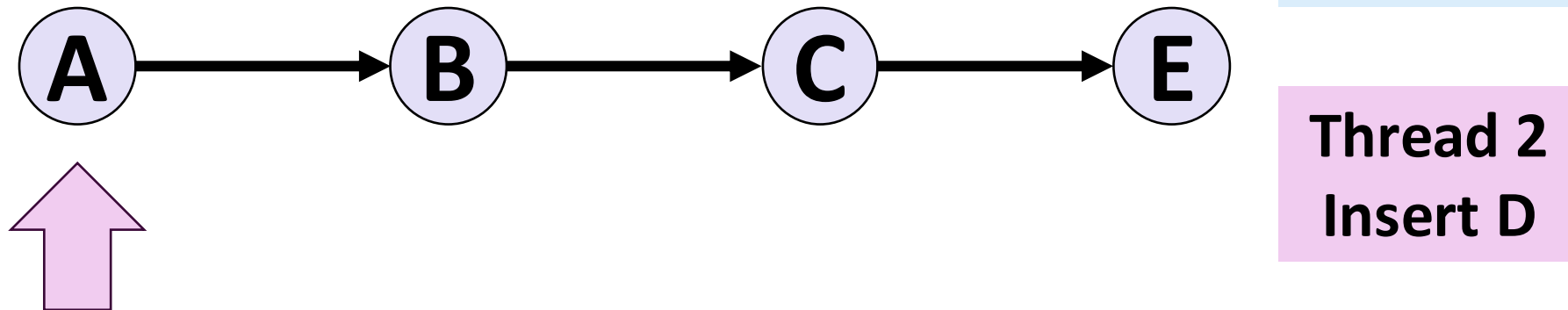
Thread 1
Delete C

Thread 2
Insert D

Optimistic Locking

- Anything else?

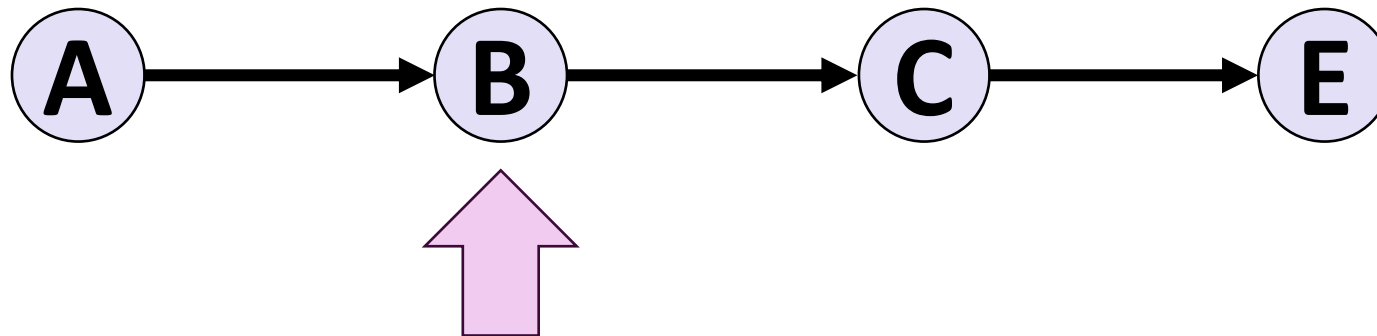
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

- What if you have a simultaneous insertion and deletion?



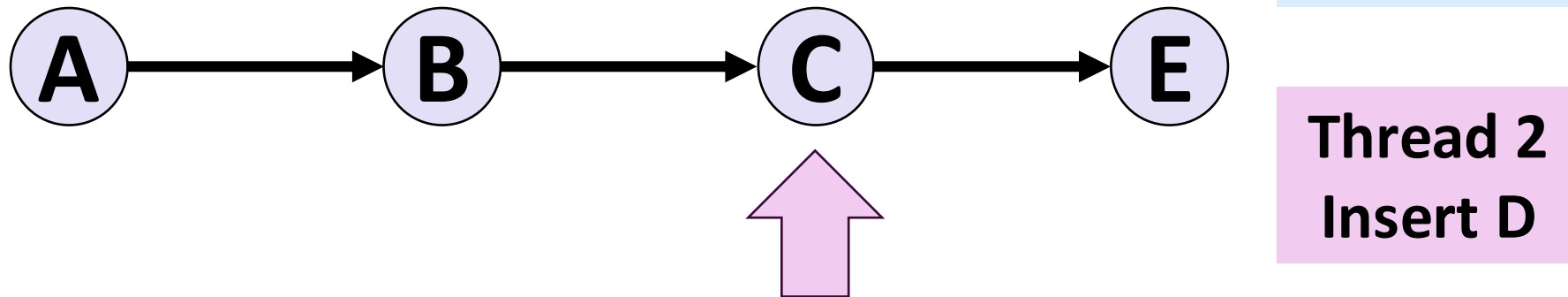
Thread 1
Delete C

Thread 2
Insert D

Optimistic Locking

- Anything else?

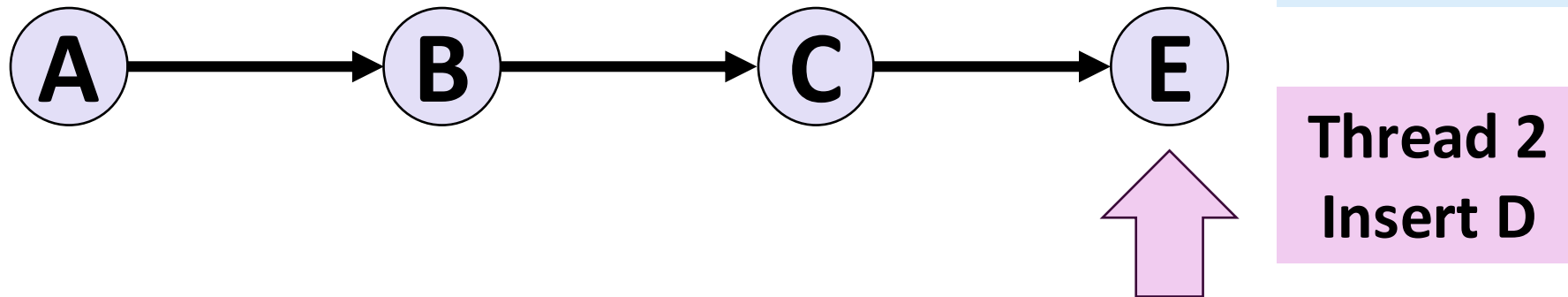
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

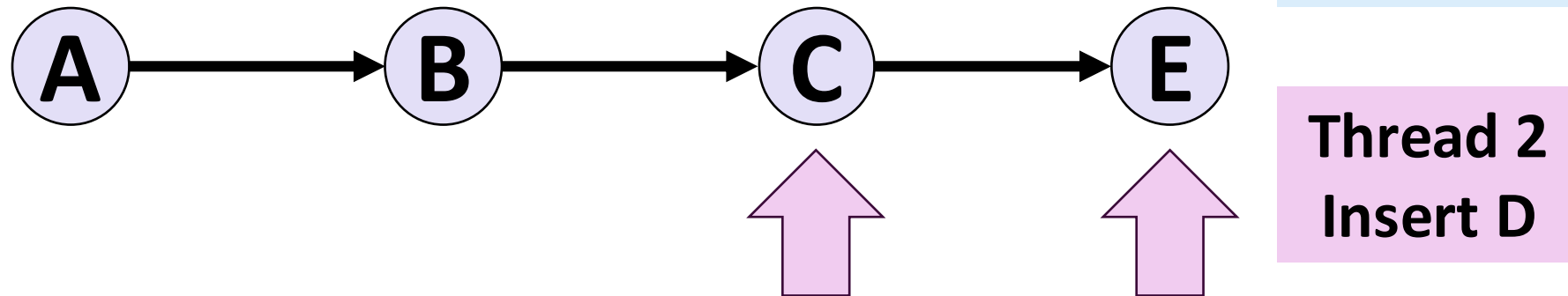
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

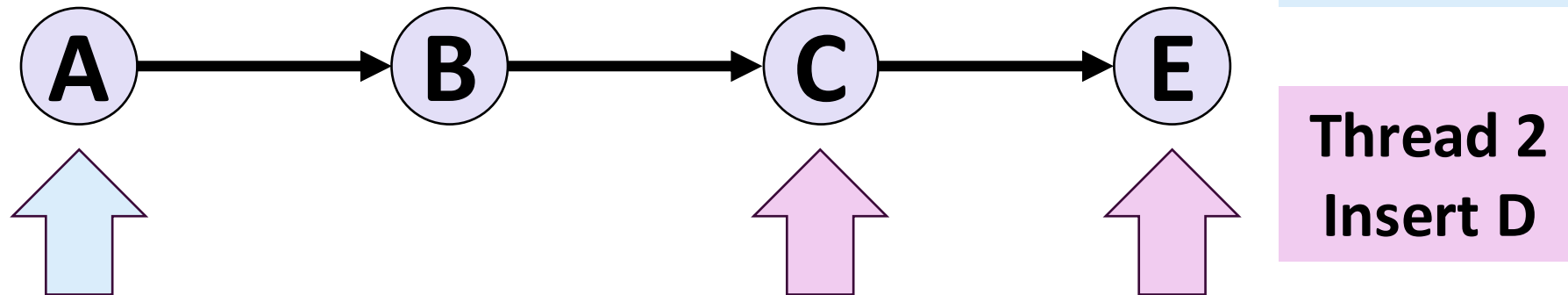
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

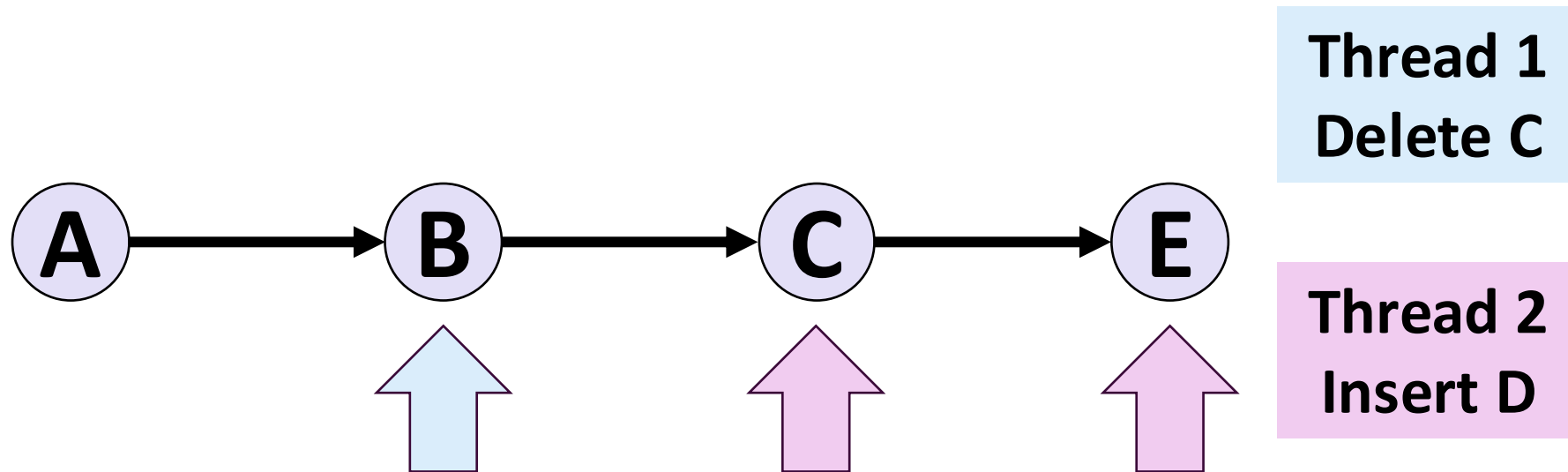
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

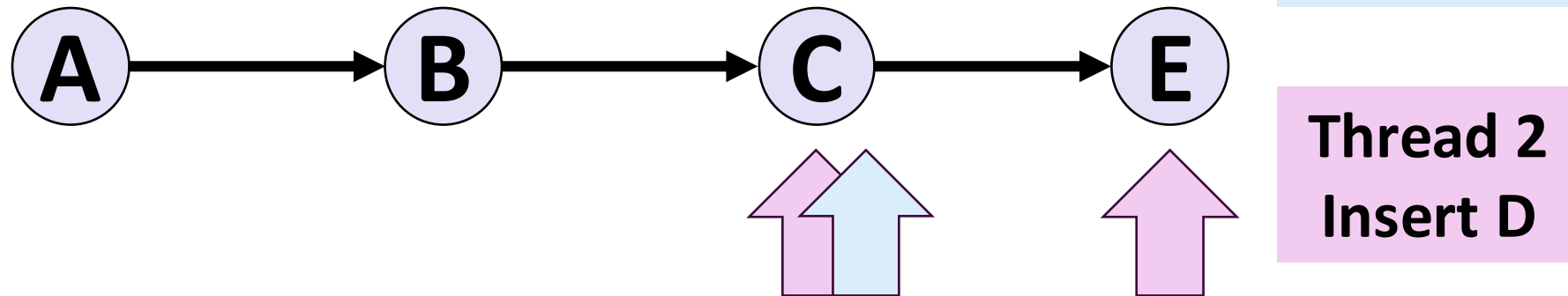
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

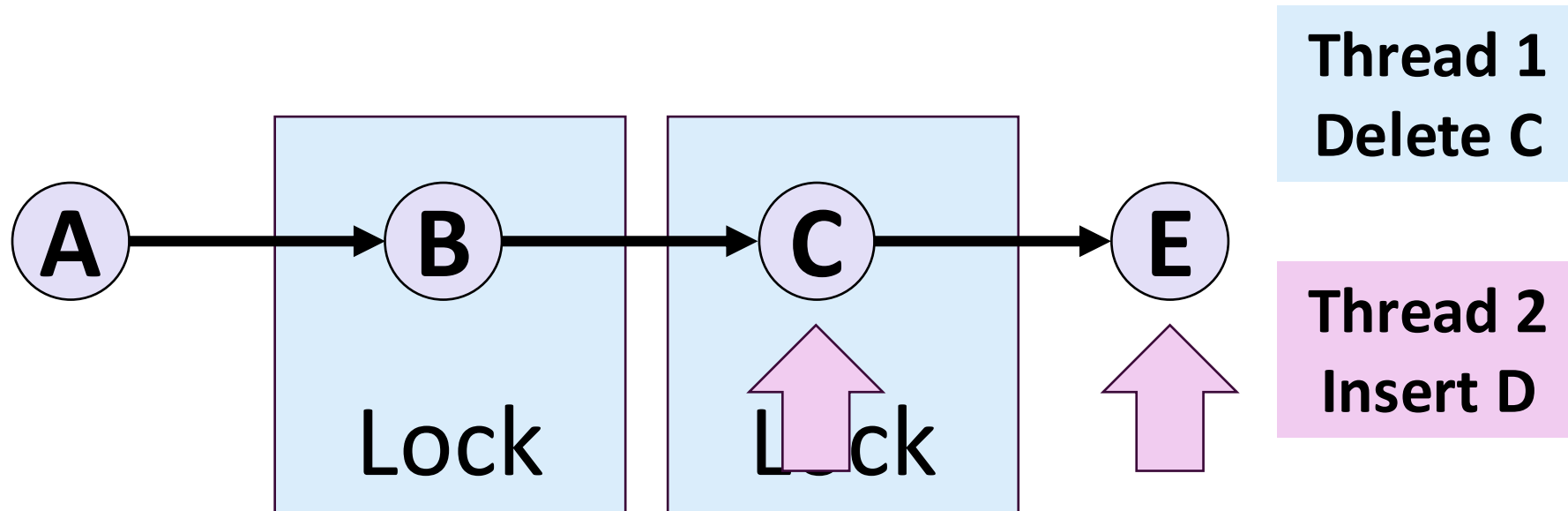
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

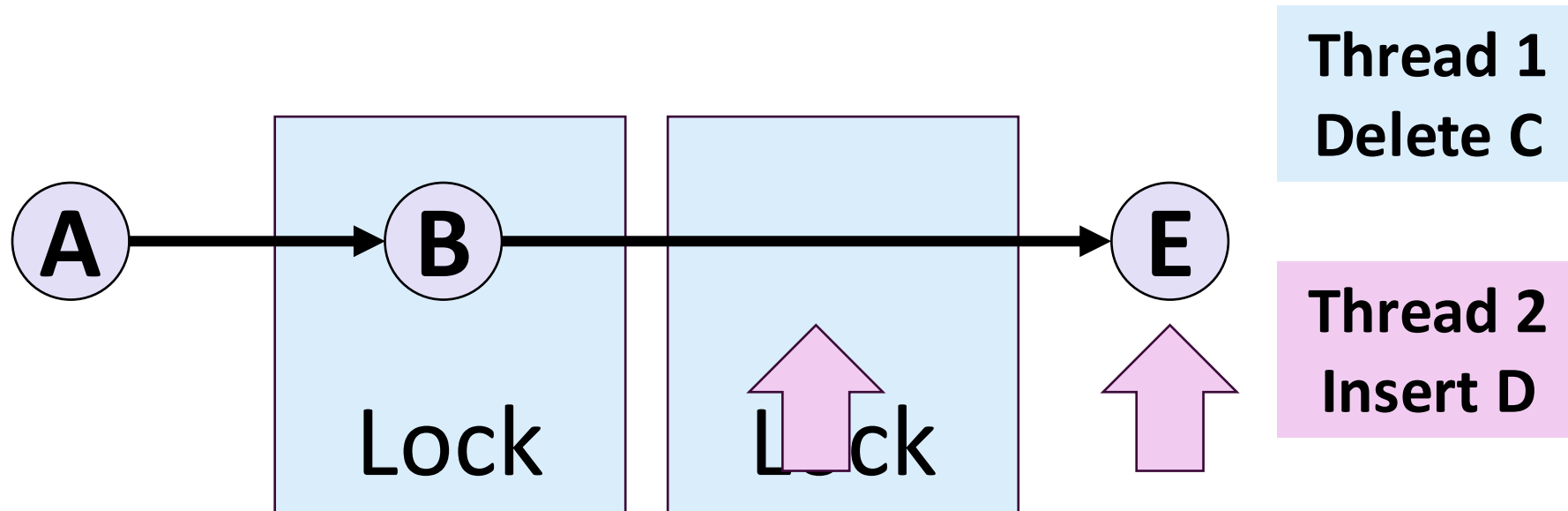
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

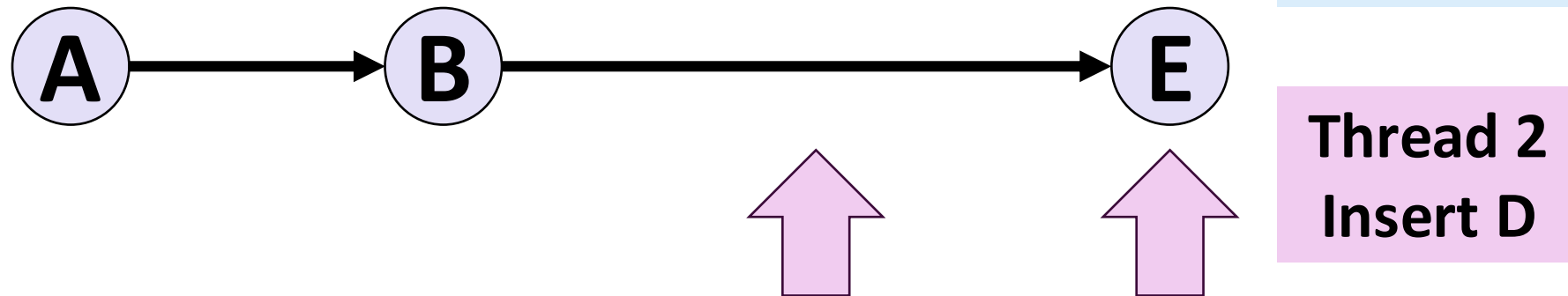
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

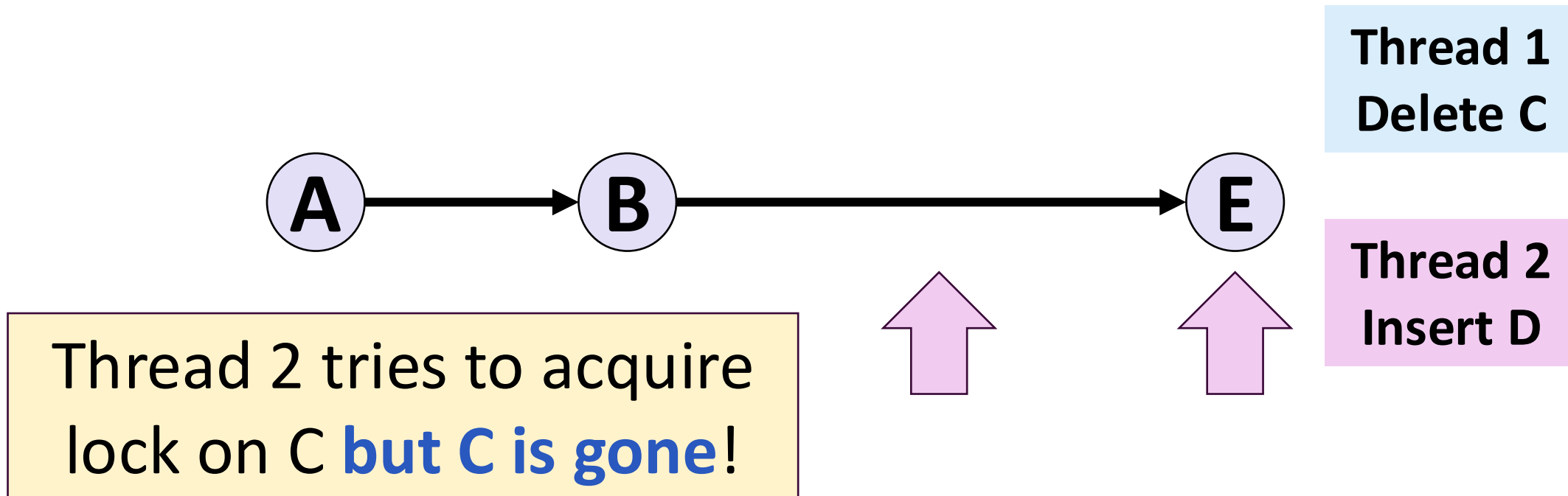
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

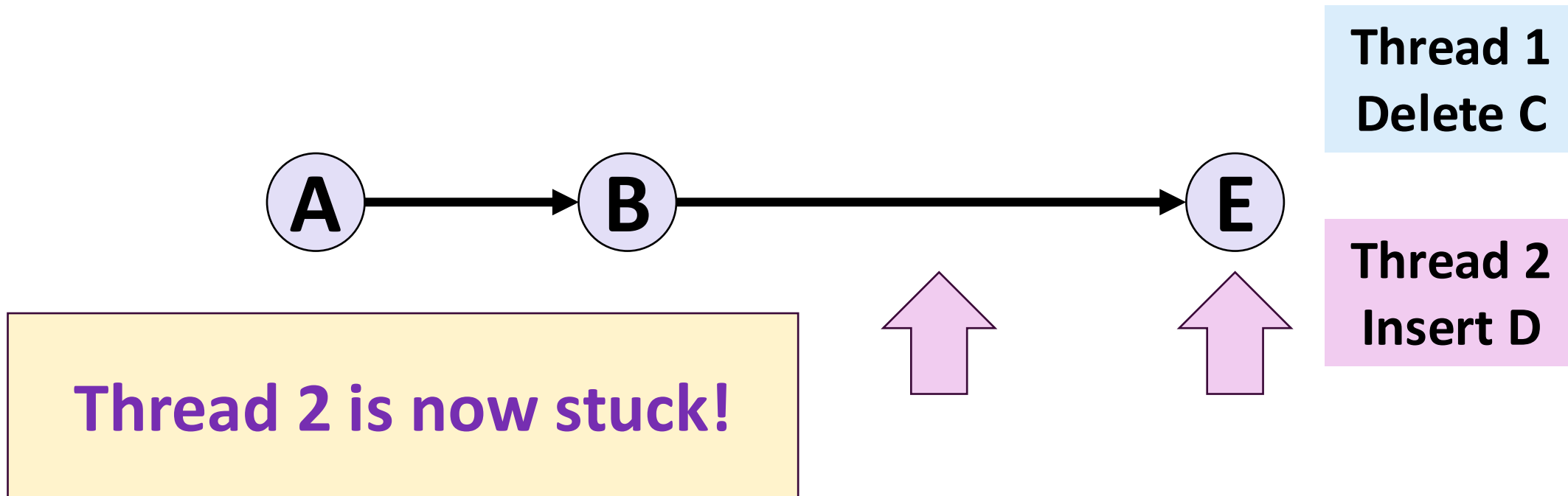
- What if you have a simultaneous insertion and deletion?



Optimistic Locking

- Anything else?

- What if you have a simultaneous insertion and deletion?



Optimistic Locking **with Validation**

- **Validation is an important concept in parallel/concurrent algorithms**
 - Before you try to acquire locks, first validate that the state of the world is what you think it is

```
bool validate(Node* pred, Node* curr) {  
    return (pred->next == curr);  
}
```

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
    }
```

```
    std::unique_lock<std::mutex> lockPred(pred->m);  
    std::unique_lock<std::mutex> lockCurr;  
    if (curr) {  
        lockCurr = std::unique_lock<std::mutex>(curr->m);  
    }
```

```
    if (validate(pred, curr)) {  
        Node* newNode = new Node(val, curr);  
        pred->next = newNode;  
        // Locks are released automatically at scope exit  
        return;  
    }  
}
```


Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
    }
```

```
    std::unique_lock<std::mutex> lockPred(pred->m);  
    std::unique_lock<std::mutex> lockCurr;  
    if (curr) {  
        lockCurr = std::unique_lock<std::mutex>(curr->m);  
    }
```

```
    if (validate(pred, curr)) {  
        Node* newNode = new Node(val, curr);  
        pred->next = newNode;  
        // Locks are released automatically at scope exit  
        return;  
    }  
}
```

Search for insertion place
without locking

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
if (validate(pred, curr)) {  
    Node* newNode = new Node(val, curr);  
    pred->next = newNode;  
    // Locks are released automatically at scope exit  
    return;  
}  
}
```

Lock only when found
insertion location

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
if (validate(pred, curr)) {  
    Node* newNode = new Node(val, curr);  
    pred->next = newNode;  
    // Locks are released automatically at scope exit  
    return;  
}  
}
```

Validate that the predecessor
and successor still contain
the link you want

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
    if (validate(pred, curr)) {  
        Node* newNode = new Node(val, curr);  
        pred->next = newNode;  
        // Locks are released automatically at scope exit  
        return;  
    }  
}
```

Add the new node

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
if (validate(pred, curr)) {  
    Node* newNode = new Node(val, curr);  
    pred->next = newNode;  
    // Locks are released automatically at scope exit  
    return;  
}
```

But there is still a problem!

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
if (validate(pred, curr)) {  
    Node* newNode = new Node(val, curr);  
    pred->next = newNode;  
    // Locks are released automatically at scope exit  
    return;  
}  
}
```

What is the problem with this code?

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
    if (validate(pred, curr)) {  
        Node* newNode = new Node(val, curr);  
        pred->next = newNode;  
        // Locks are released automatically at scope exit  
        return;  
    }  
}
```

Node pred could be deleted before you attempt to lock it

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
if (validate(pred, curr)) {  
    Node* newNode = new Node(val, curr);  
    pred->next = newNode;  
    // Locks are released automatically at scope exit  
    return;  
}  
}
```

**One fix: do not delete nodes
from memory; only delete
the pointers**

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next;  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
  
        std::unique_lock<std::mutex> lockPred(pred->m);  
        std::unique_lock<std::mutex> lockCurr;  
        if (curr) {  
            lockCurr = std::unique_lock<std::mutex>(curr->m);  
        }  
    }  
}
```

```
    if (validate(pred, curr)) {  
        Node* newNode = new Node(val, curr);  
        pred->next = newNode;  
        // Locks are released automatically at scope exit  
        return;  
    }  
}
```

However, this could result in unbounded memory!

Optimistic Locking with Validation

```
void insert(int val) {  
    while (true) {  
        Node* pred = head;  
        Node* curr = pred->next  
        while (curr && curr->value < val) {  
            pred = curr;  
            curr = curr->next;  
        }  
    }
```

```
    std::unique_lock<std::mutex> lockPred(pred->m);  
    std::unique_lock<std::mutex> lockCurr;  
    if (curr) {  
        lockCurr = std::unique_lock<std::mutex>(curr->m);  
    }
```

```
    if (validate(pred, curr)) {  
        Node* newNode = new Node(val, curr);  
        pred->next = newNode;  
        // Locks are released automatically at scope exit  
        return;  
    }  
}
```

How do we fix this new problem? Think about it for next class.

Runtimes of different locking mechanisms

Linked List with 500 elements

Hand-over-Hand Locking: 39116 ms

Optimistic Locking: 6447 ms

6x improvement!

Lecture 10 Quiz

